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Minimisation of Event Structures



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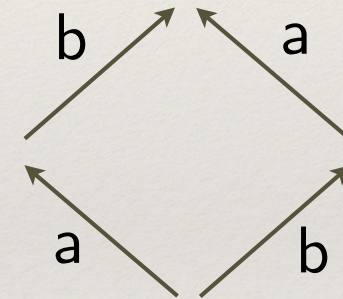
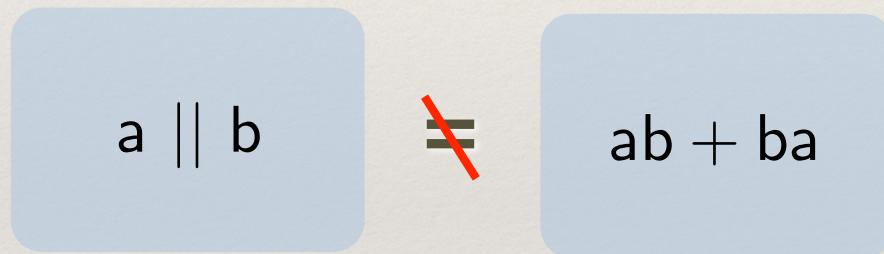
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What are we talking about?

True Concurrency

❖ Interleaving

concurrency reduced to the nondeterministic choice between the possible interleaving



True concurrent
semantics

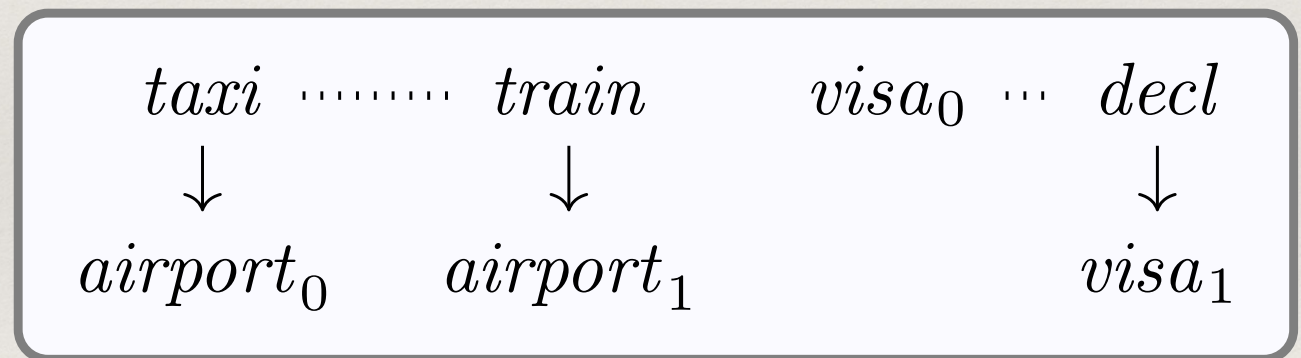
different
concurrency
causality

(Prime) Event structures

❖ **Prime event structures** [Nielsen, Plotkin, Winskel]

$(E, \leq, \#, \lambda)$

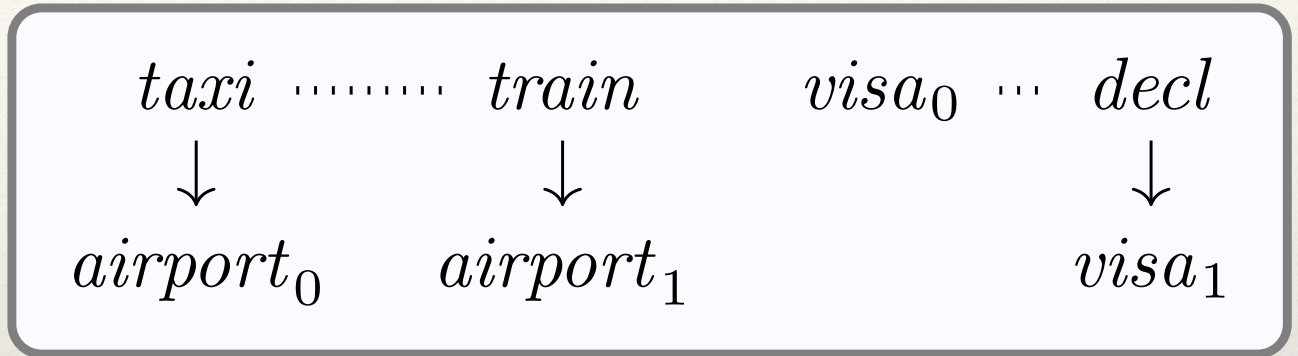
- E events
- \leq causality
- $\#$ conflict
- λ labelling



Computations as Configurations

$(E, \leq, \#, \lambda)$

Computations as
configurations



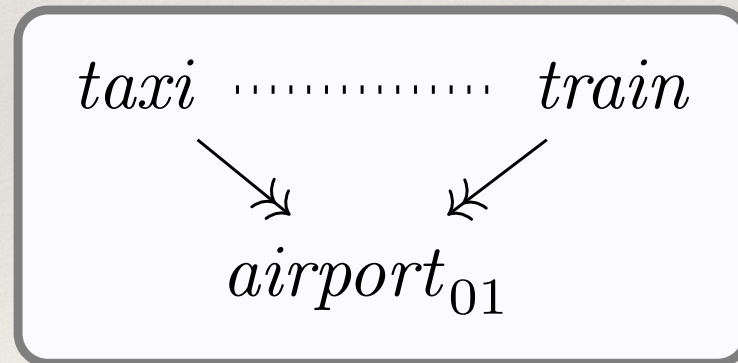
(causally closed, conflict-free)

$$\emptyset \xrightarrow{taxi} \{taxi\} \xrightarrow{visa_0} \{taxi, visa_0\} \xrightarrow{airport_0} \{taxi, visa_0, airport_0\}$$

More expressive ES models

❖ **Flow event structures** [Boudol, Castellani]

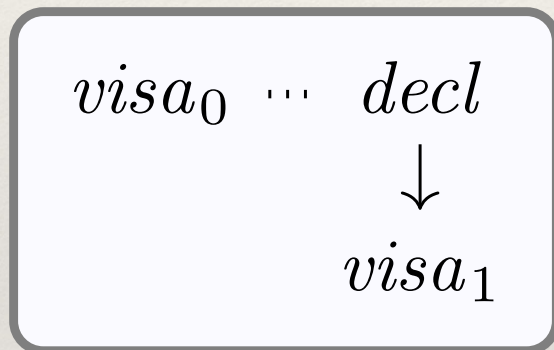
- Causality replaced by **flow relation** which permits conflictual disjunctive causes



More expressive ES models

❖ **Asymmetric event structures**

[Langerak], [B.. Corradini, Montanari]

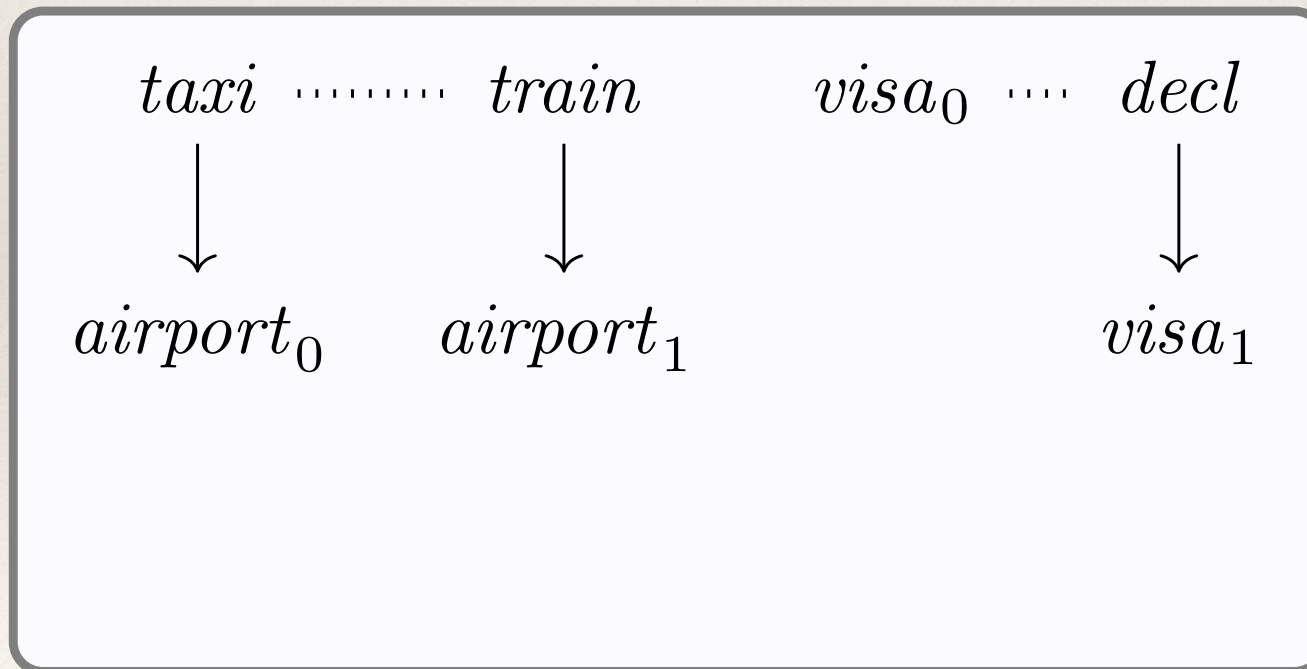


- *Asymmetric conflict:*
 $decl$ can be executed only before $visa_{01}$

Mining processes as ESs

- ❖ Processes mined in the form of event structures

[Dumas, García-Bañuelos, Armas-Cervantes, ...]

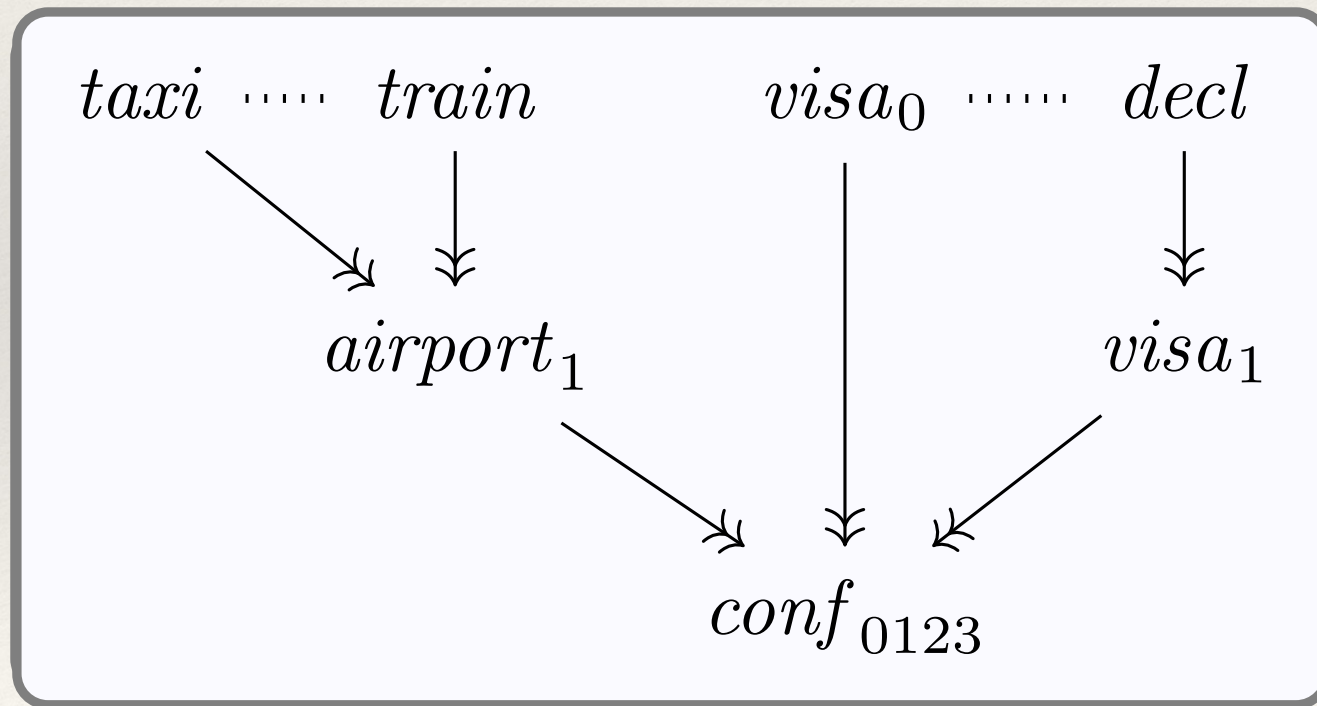


Original questions

- ❖ Can we get **smaller models**, by some form of **quotient**, keeping the **concurrent behaviour unchanged**?
- ❖ **Folding**: Surjective $f : E \rightarrow E'$, that **merges some events** establishing a **concurrent bisimulation**
- ❖ Some work ad hoc in [B., García-Bañuelos, Armas-Cervantes,]

Original questions

❖ It worked ...



General theory?

General theory of behaviour-preserving quotients for event structures?

- ❖ Is the **notion of folding adequate** for expressing behaviour-preserving quotients?
- ❖ Is there a **minimal quotient** in some general class of event structures? Does it exist in **specific subclasses**?
- ❖ Can we have a **characterisation** of foldings directly on the ES?

Poset ES

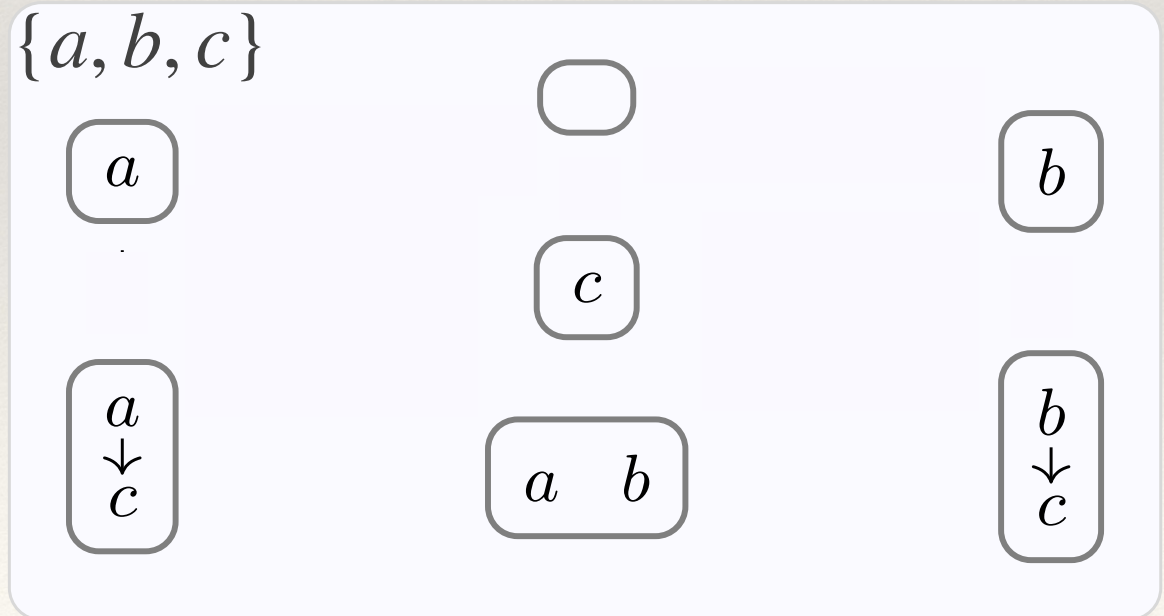
❖ $\langle E, Conf(E), \lambda \rangle$

[Rensink]

❖ E events

❖ $Conf(E)$ is a set of **configurations** (posets of events)

❖ **Prefix order**
on configurations



Stable ES as Poset ES

- ❖ **Common stable ES models are instances of poset ES**
 - ❖ prime ES [Nielsen, Plotkin, Winskel]
 - ❖ flow ES [Boudol, Castellani]
 - ❖ bundle ES [Langerak]
 - ❖ asymmetric ES. [B., Corradini, Montanari]
 - ❖ ...

Behavioural Equiv.: Hhp-bisimilarity

fully observes the interplay between
causality, concurrency, non-determinism

HHP-bisimilarity between E and E'

$$R = \{(C, f, C') \mid C \in \text{Conf}(E), C' \in \text{Conf}(E'), f : C \xrightarrow{\sim} C'\}$$

- Simulation:

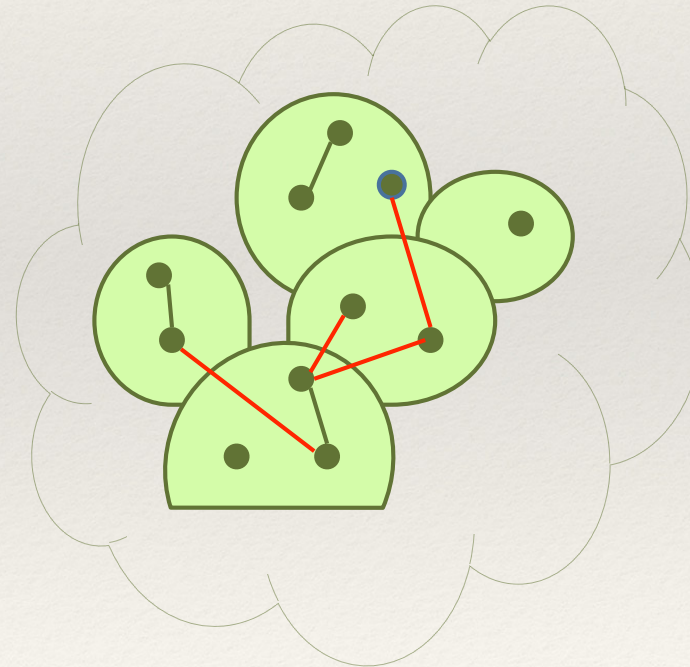
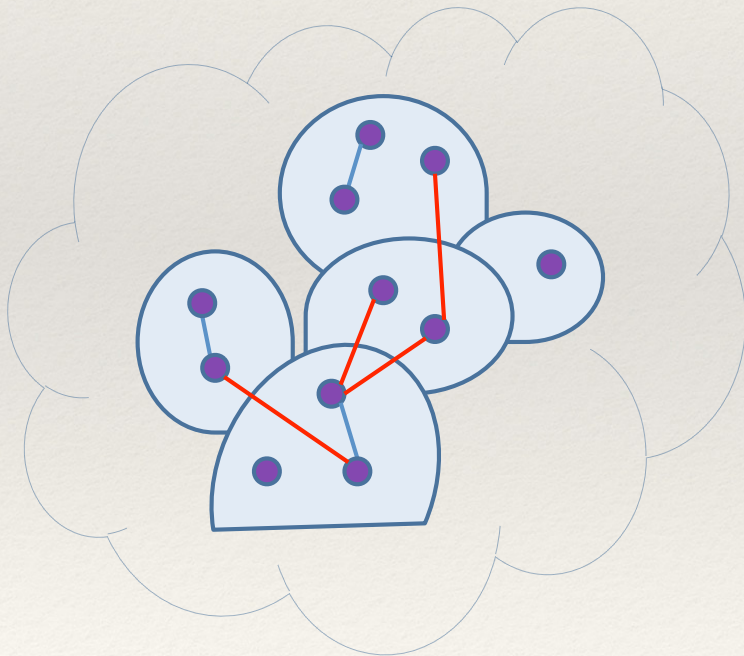
if $(C_1, f, c'_1) \in R$ and $C_1 \xrightarrow{x} C_2$ then $C'_1 \xrightarrow{x'} C'_2$
with $(C_2, f[x \mapsto x'], C'_2) \in R$, and vice versa.

- Down closure:

if $(C_1, f, c'_1) \in R$ and $C_2 \sqsubseteq C_1$ then $(C_2, f|_{C_2}, f(C_2)) \in R$.

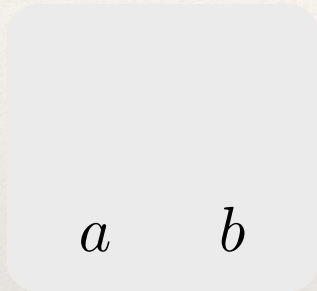
HHP-Bisimilarity

- ❖ An event of a system must be simulated by an event of the other with the same history (causal links)



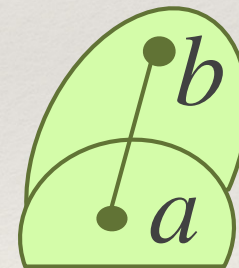
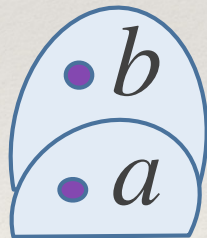
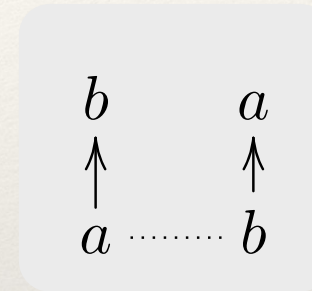
Example

$E_1: a \parallel b$

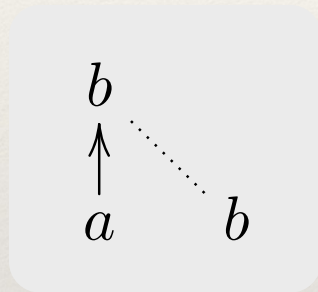


$\mathcal{F}hhp$

$E_2: ab + ba$

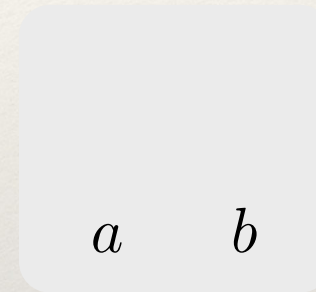


E₁



$\not\sim_{hhp}$

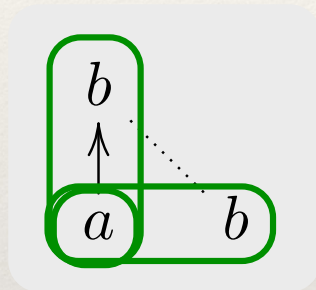
E₂



$$\emptyset \xrightarrow{\{a\}} \{a\} \xrightarrow{\{b\}} \{a < b\}$$

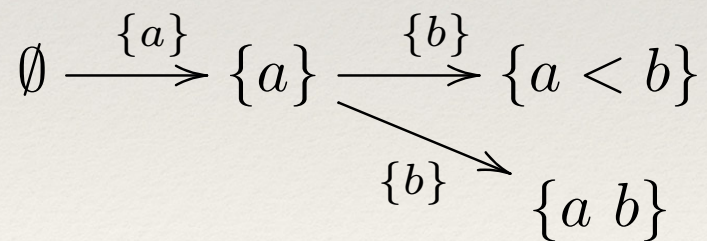
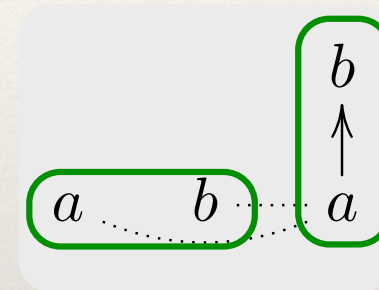
E₁ can perform *b* **causally dependent** on *a* while
E₂ cannot

E_1



$\not\sim_{hhp}$

E_2



after a , E_1 can choose between a **causally dependent** and a **concurrent** b while E_2 cannot

Folding

Morphism

$f: E_1 \rightarrow E_2$ function on events, transforms confs into confs

ES

Folding

Surjective morphism $f: E \rightarrow E'$ such that

$$R_f = \{(C, f|_C, f(C)) \mid C \text{ is a conf}\}$$

is hhp-bisimilarity

ES_f

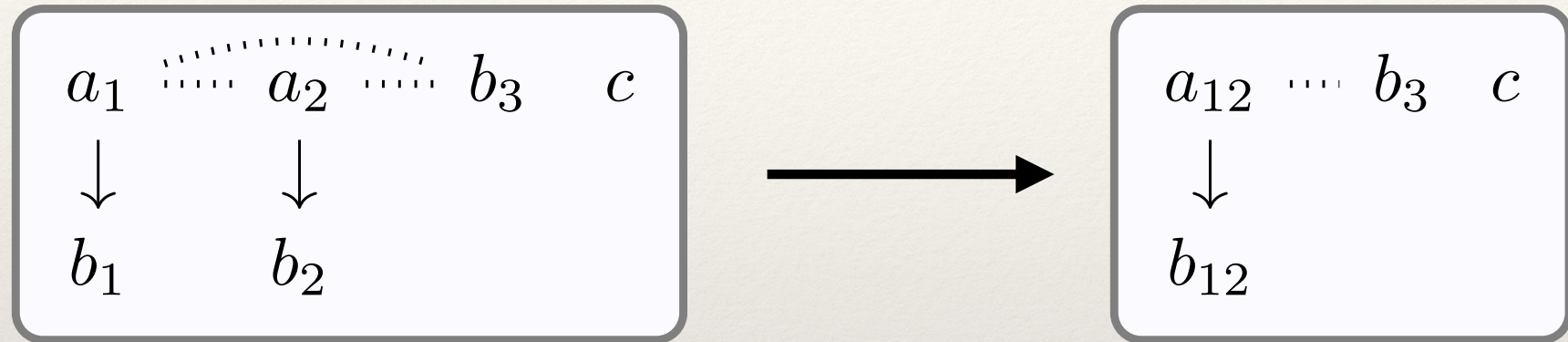
Int.: Merge events without changing the concurrent behaviour

Folding Equivalence \equiv_f

def. by $e_1 \equiv_f e_2$ iff $f(e_1) = f(e_2)$

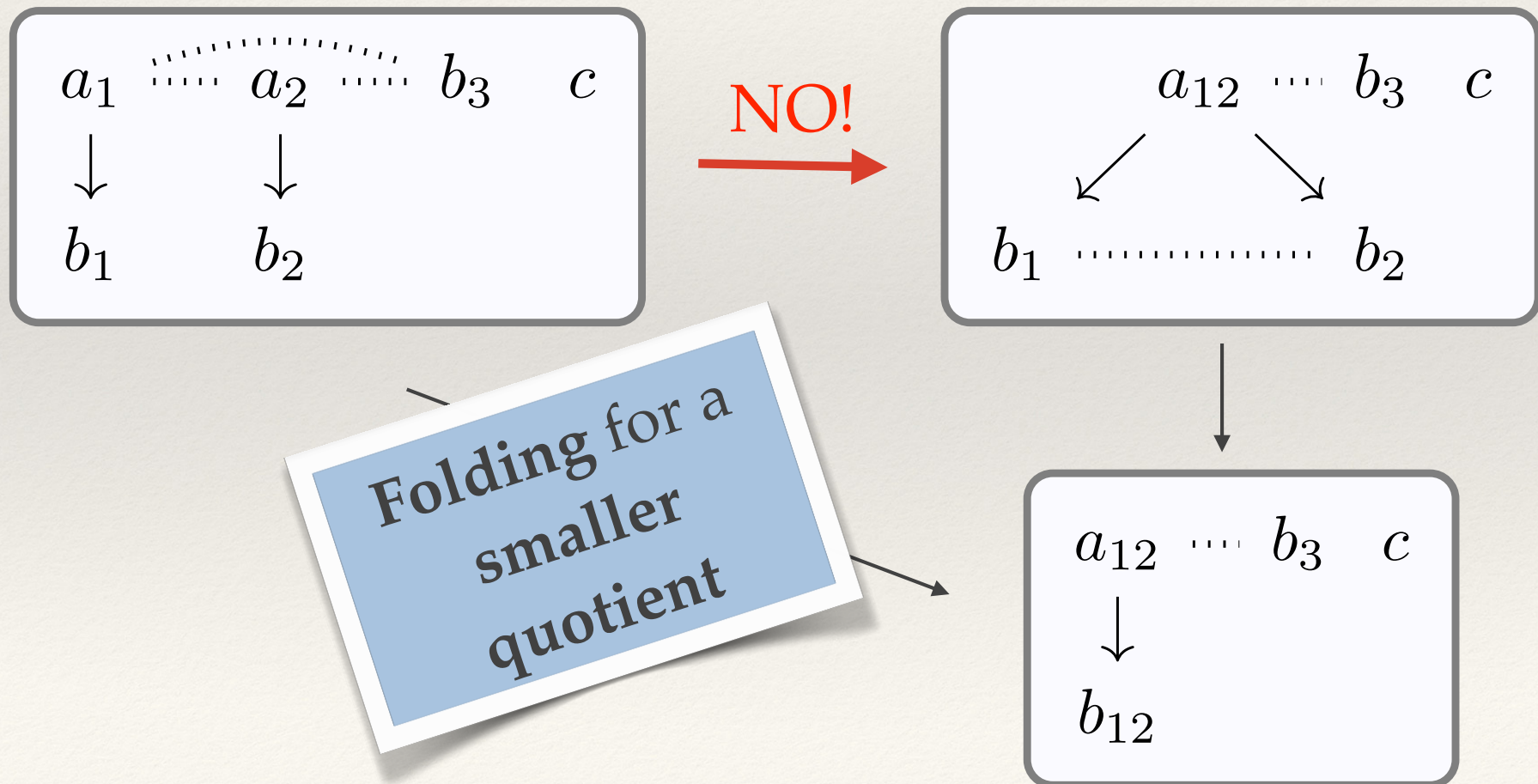
then $E \sim_{hhp} E_{/\equiv_f}$

Folding: Example



Is the notion of folding adequate?

- ❖ Not all quotients which preserve behaviour are foldings

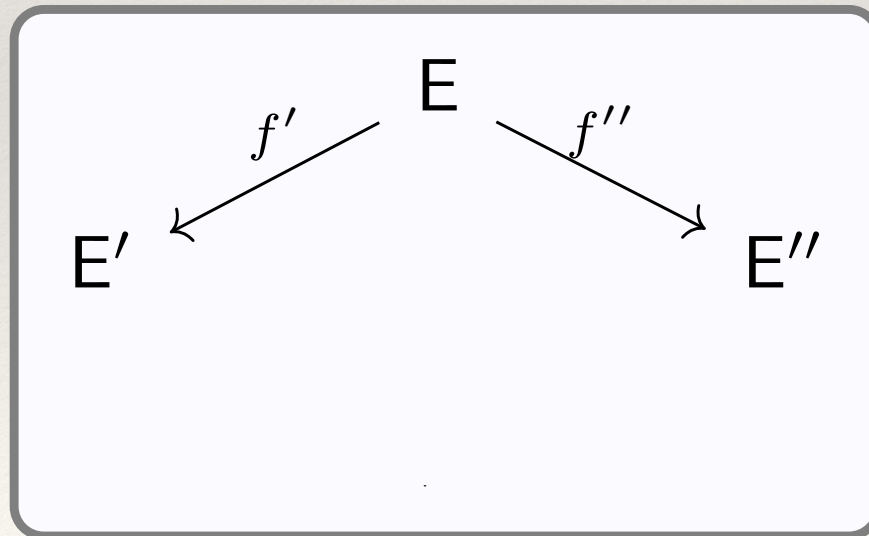


Minimal quotients?

Joining foldings

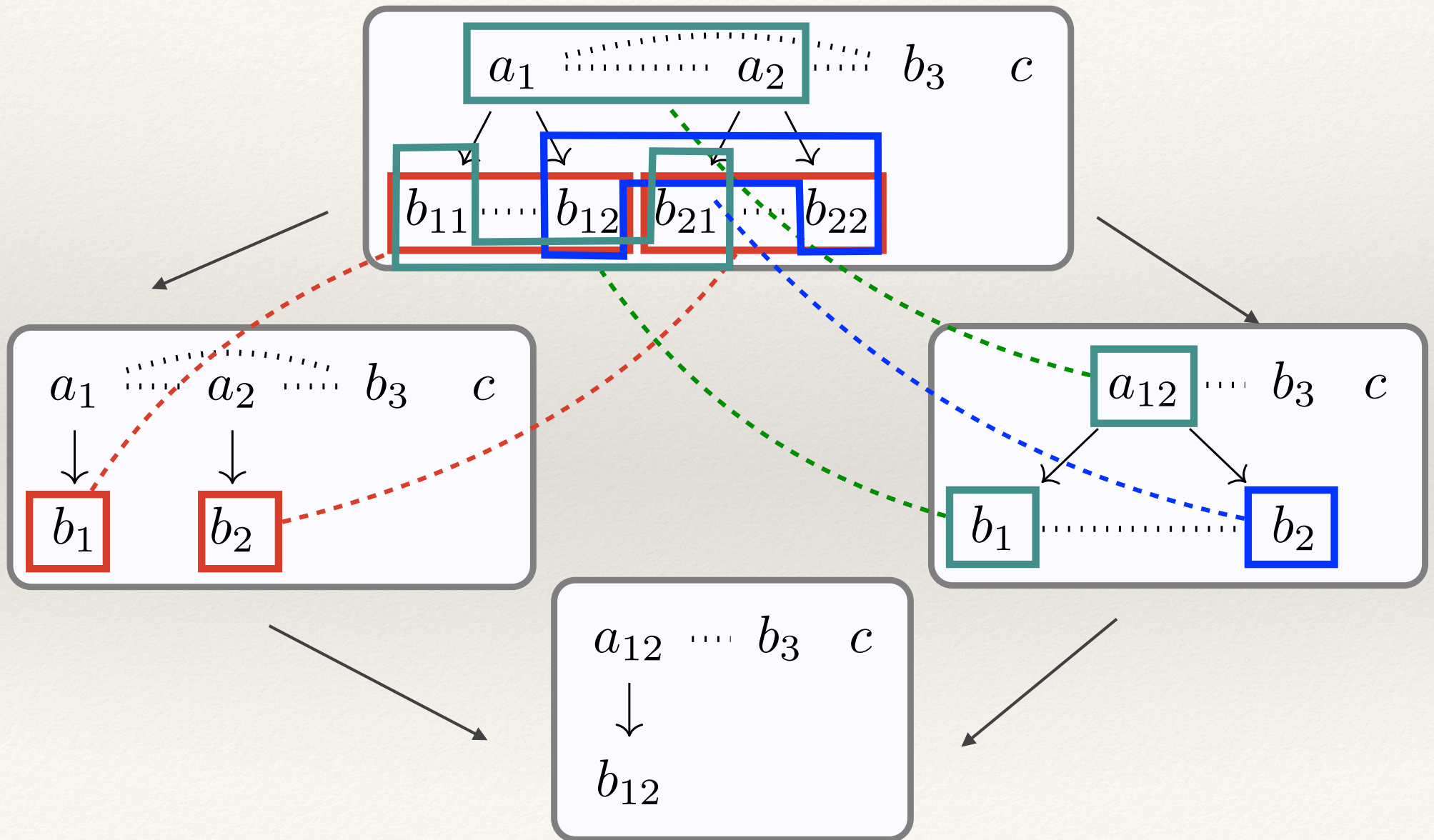
Joining Foldings

Given two foldings $f : E \rightarrow E'$ and $g : E \rightarrow E''$ there are always foldings $f : E' \rightarrow E_{/\equiv}$ and $g : E'' \rightarrow E_{/\equiv}$ where \equiv is $(\equiv_f \cup \equiv_g)^*$



Pushout in
 ES_f

Joining foldings: example



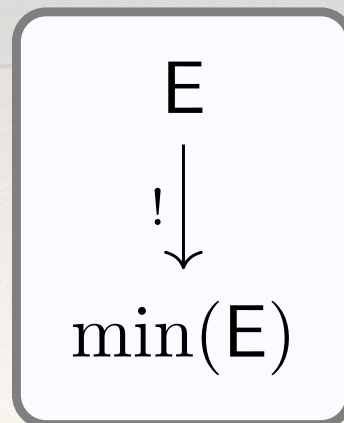
Maximally folded ES / Minimal quotient

Lattice of folding equivalences

Given a poset ES the **folding equivalences** form a complete sublattice of the lattice of equivalences over E

Corollary

Given a poset ES there exists its maximally folded version $\min(E)$

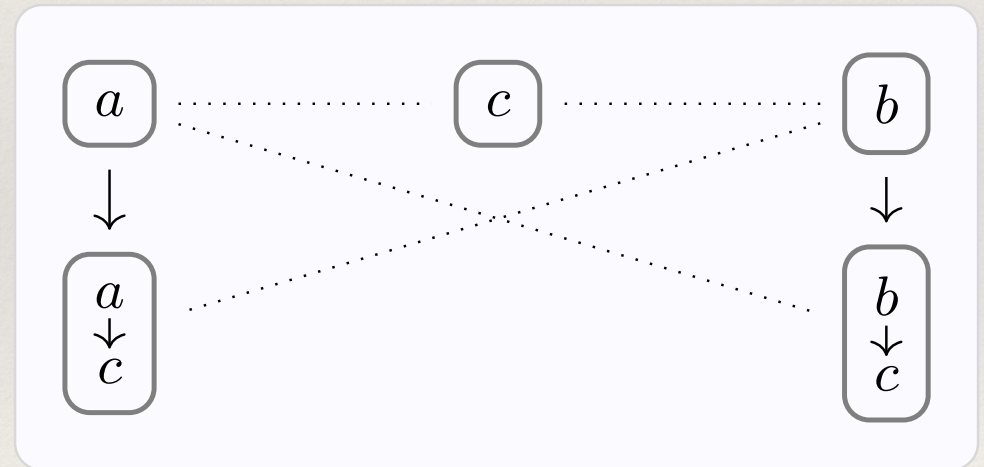
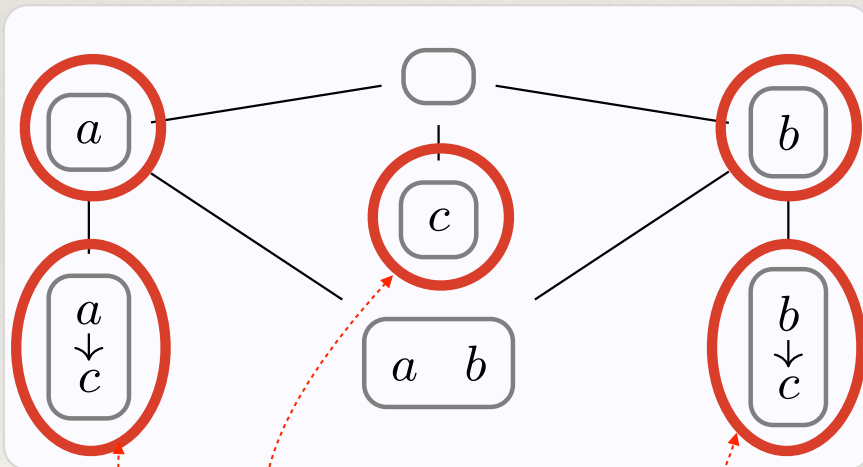


Characterising foldings

Folding through PESs

Theorem: Each poset ES E is the folding of a canonical Prime ES

$$\phi_E : \mathbb{P}(E) \rightarrow E$$



Histories of event c

Folding factorises through PESs

Theorem

Given poset ES E and E' and a morphism $f : E \rightarrow E'$
consider

$$E \xrightarrow{f} E'$$

f is a folding iff $\mathbb{P}(f)$ is a folding

Foldings on PES

Characterisation

$f : P \rightarrow P'$ surjective PES isomorphism

f is a folding iff $\forall W \subseteq P$ $f(W) \in P'$

1. if $x \#^{\forall} f^{-1}(y)$ then $f(x) \# y'$;

2. if $f(x) = f(y)$, $\bigcap W \ni x, y$ then $\exists z \in P$ such that $f(z) = f(x)$ and $\bigcap (W \cup \{z\})$.

Labels & causes

Conflicts

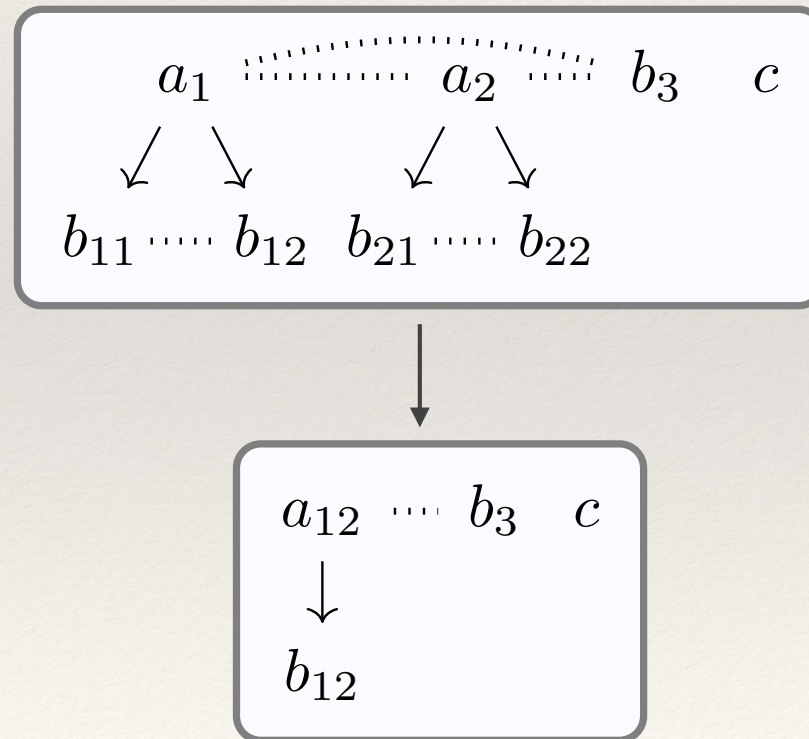
Concurrency

Minimal Foldings for PESs?

Theorem:

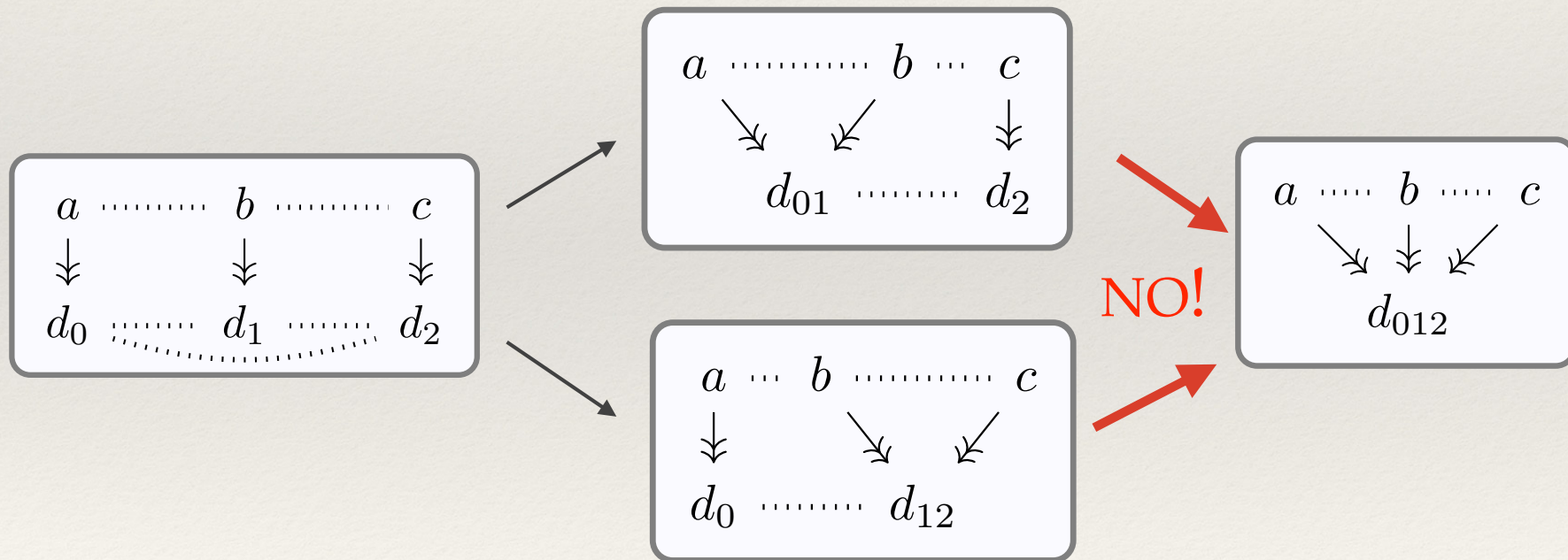
Any PES has maximally folded version in the class of PESs

Example



Minimal Foldings for other classes?

Not true in the class of FESs (and other subclasses)



Conclusion

Foldings as **open maps** (with poset as cat of obs) [Joyal, Nielsen, Winskel]
-> more abstract and general view of the results?

Folding algorithms (for PESs and other subclasses)?

Folding for operational models (finitary representations of ESs, e.g. Petri nets)