

A Diagrammatic Algebra of Linear and Concurrent Systems

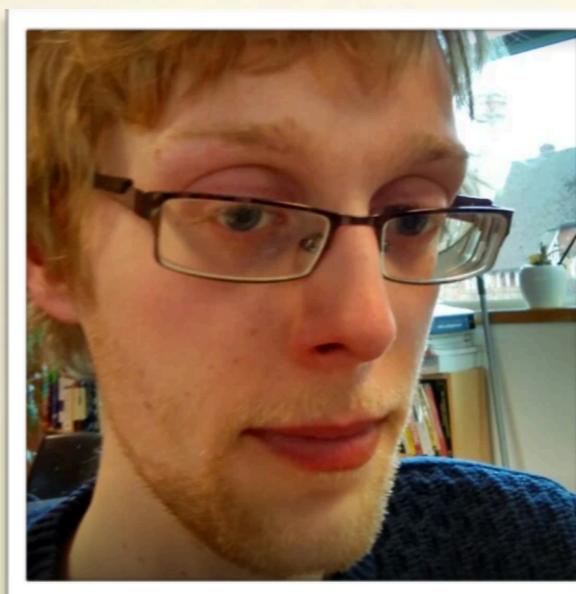
Fabio Zanasi
University College London

IFIP WP1.3 Meeting, Prague
April 2019

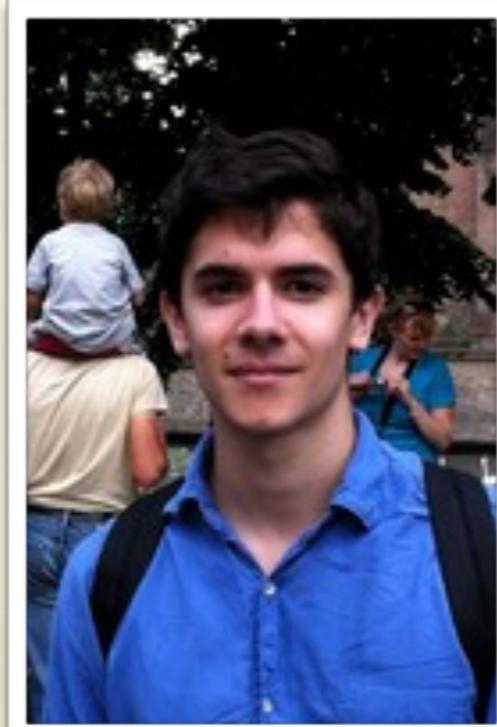
Collaborators



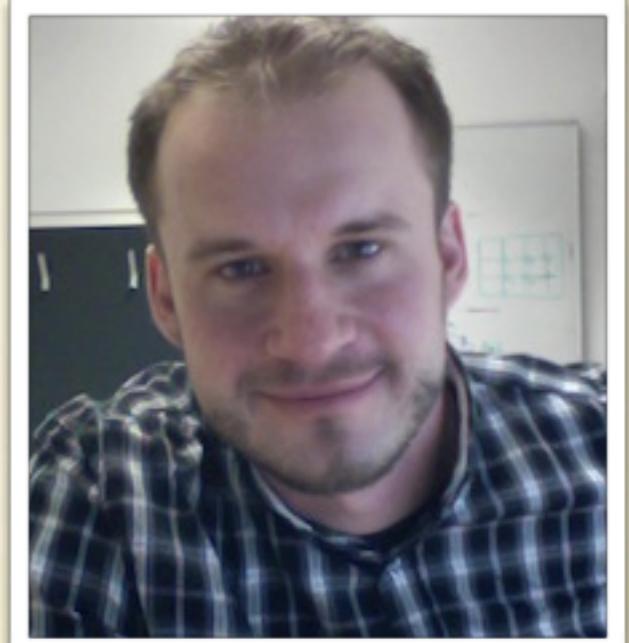
Filippo
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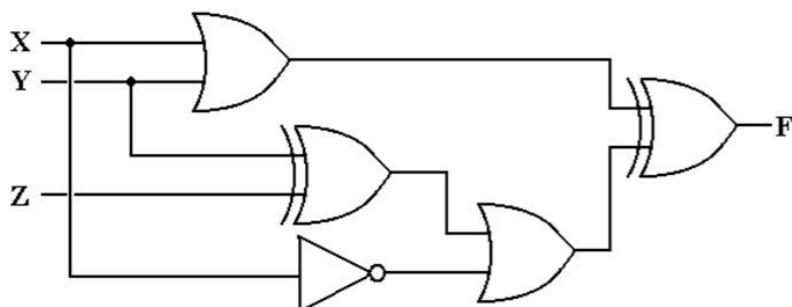


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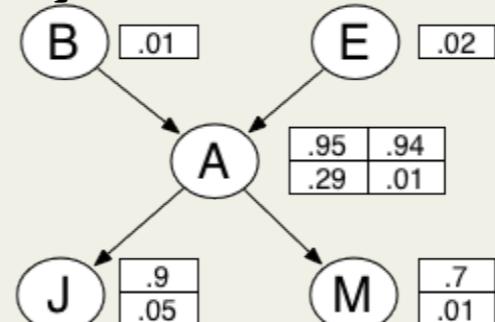
Introduction

Component-Based Systems

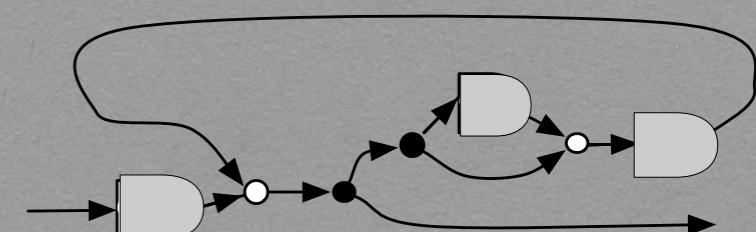
Digital Circuits



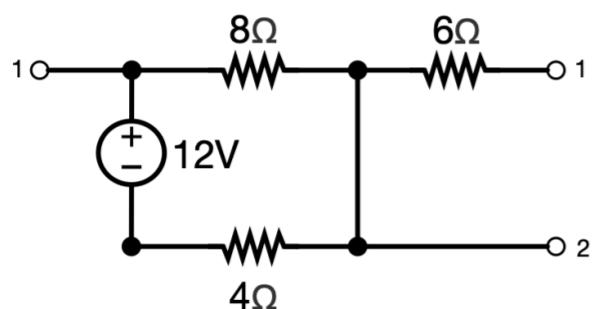
Bayesian Networks



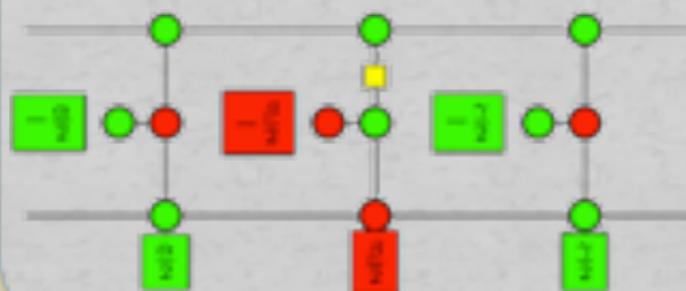
Signal Flow Graphs



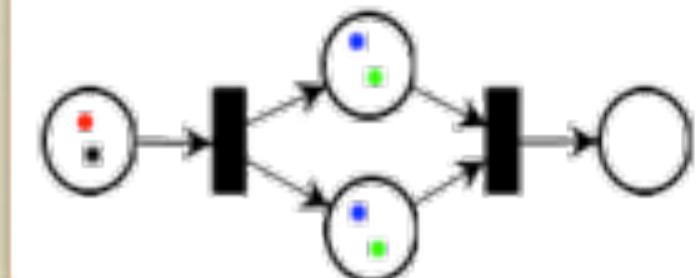
Electrical Circuits



Quantum Processes

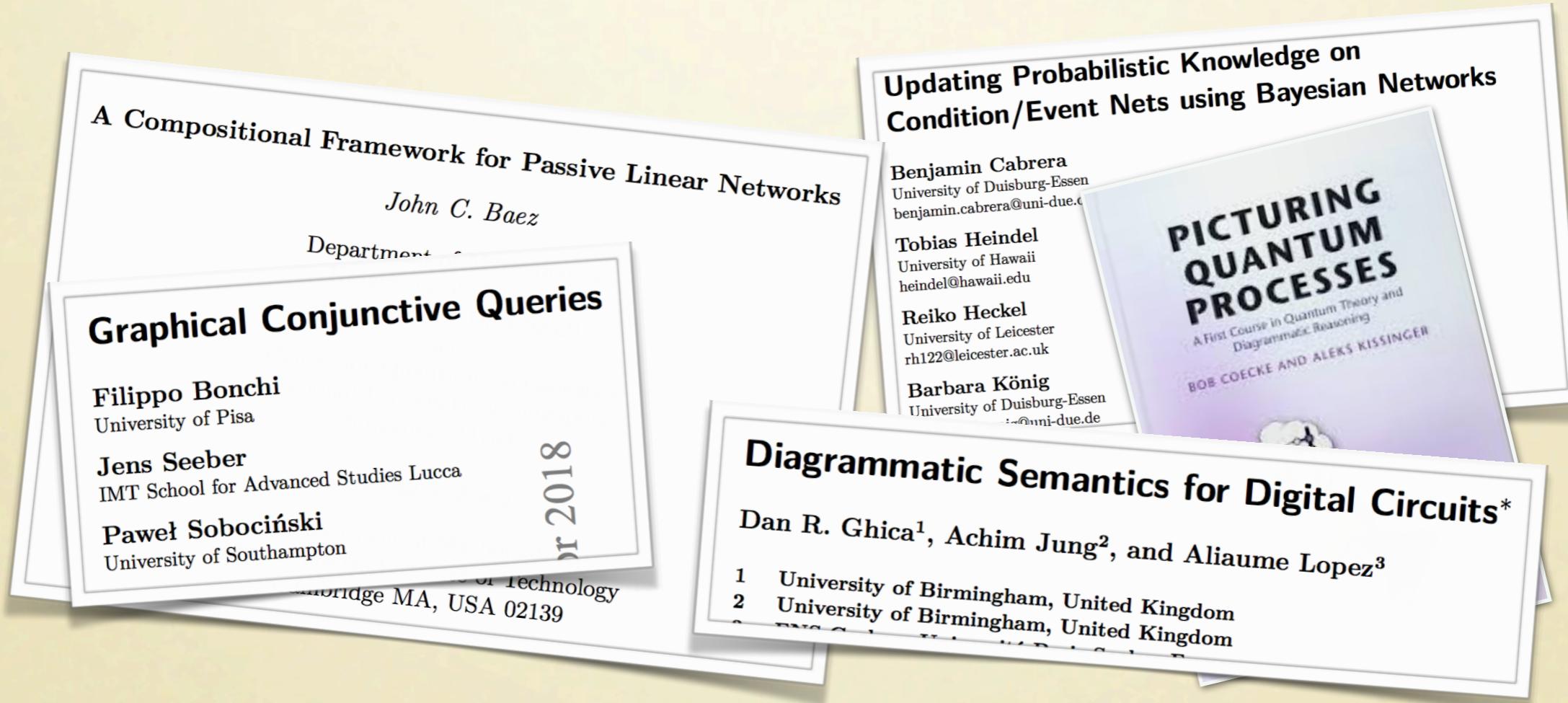


Petri Nets



Compositional Modelling

There is an emerging, multi-disciplinary field aiming at studying graphical models of computation **compositionally**, inspired by the **algebraic methods** of programming language semantics.



Diagrams are first-class citizens of the theory. The appropriate algebraic setting is **monoidal** (and not **cartesian**) categories.

In this talk

Signal flow interpretation:
linear relations

Resource interpretation:
additive relations

Fully complete
axiomatisation

Fully complete
axiomatisation

Diagrammatic
circuit syntax
GLA

Signal Flow Graphs

Petri nets

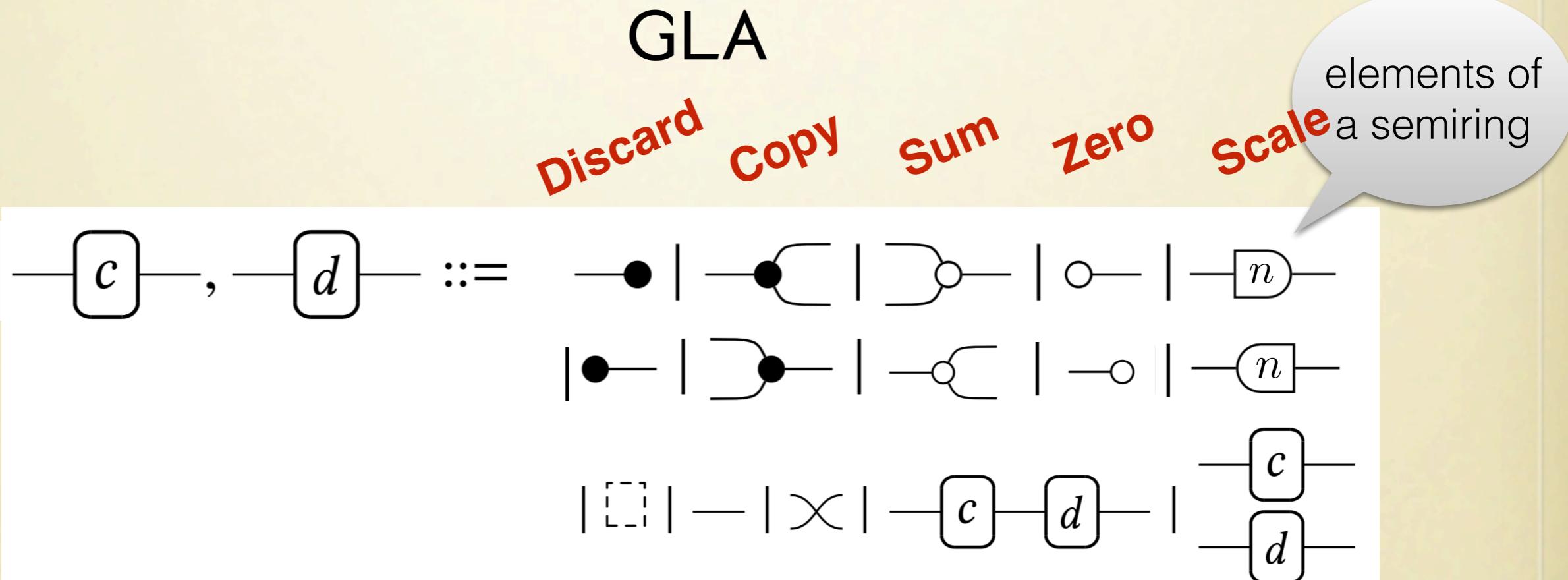
Take home message:
seemingly diverse computational models can be studied
within **the same** algebraic framework.

The core language

A simple diagrammatic syntax

Graphical Linear Algebra

GLA



Circuit Diagrams Behaviour

$$\bullet \xrightarrow[n]{\cdot} \bullet \quad \bullet \xrightarrow[n]{n} \bullet \quad \bullet \xrightarrow[n]{n} \bullet$$

$$\circlearrowleft \xrightarrow[m]{n+m} \circlearrowright \quad \circlearrowleft \xrightarrow[0]{\cdot} \circlearrowright$$

$$\square[n] \xrightarrow[mn]{m} \square[n]$$

$$\bullet \xrightarrow[n]{\cdot} \bullet \quad \bullet \xrightarrow[n]{n} \bullet \quad \bullet \xrightarrow[n]{n} \bullet$$

$$\circlearrowleft \xrightarrow[n]{n+m} \circlearrowright \quad \circlearrowleft \xrightarrow[0]{\cdot} \circlearrowright$$

$$\square[n] \xrightarrow[mn]{m} \square[n]$$

$$[] \xrightarrow[\cdot]{\cdot} [] \quad - \xrightarrow[n]{n} - \quad \times \xrightarrow[m]{m} \times$$

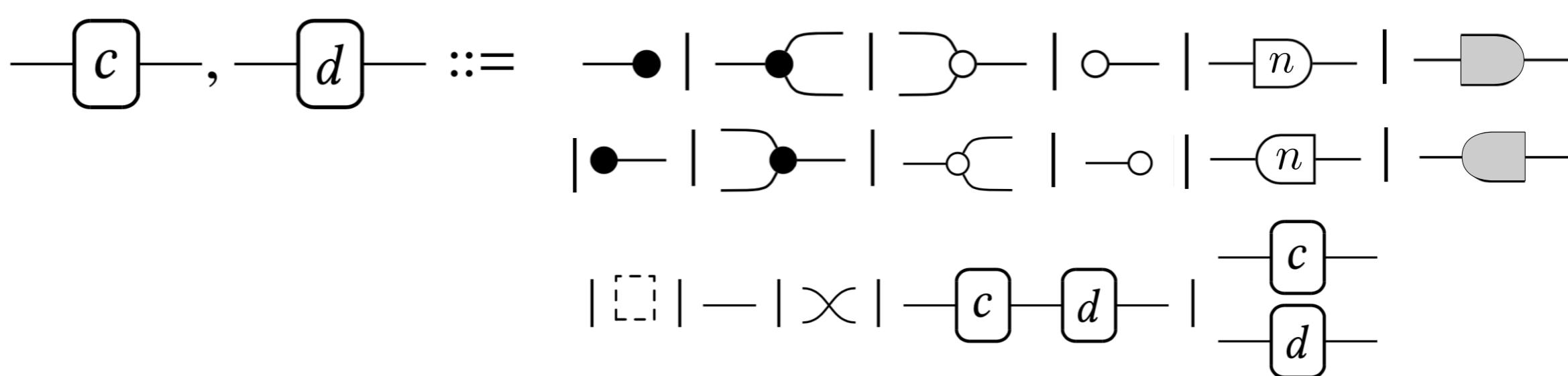
$$\begin{array}{c} \boxed{c_1} \xrightarrow[b]{a} \boxed{c_2} \quad \boxed{c_3} \xrightarrow[c]{b} \boxed{c_4} \\ \hline \boxed{c_1} \boxed{c_3} \xrightarrow[c]{a} \boxed{c_2} \boxed{c_4} \end{array}$$

$$\begin{array}{c} \boxed{c_1} \xrightarrow[b_1]{a_1} \boxed{c_2} \quad \boxed{c_3} \xrightarrow[b_2]{a_2} \boxed{c_4} \\ \hline \boxed{c_1} \boxed{c_3} \xrightarrow[b_1]{a_1} \boxed{c_2} \boxed{c_4} \xrightarrow[b_2]{a_2} \end{array}$$

$$[[c]] := \{(a, b) \mid c \xrightarrow[b]{a} c\}$$

Stateful extension: syntax

GLA_s



Stateful extension: behaviour

$$(\text{--} \square \text{--}, m) \xrightarrow[m]{n} (\text{--} \square \text{--}, n)$$

$$(\text{--} \square \text{--}, m) \xrightarrow[n]{m} (\text{--} \square \text{--}, n)$$

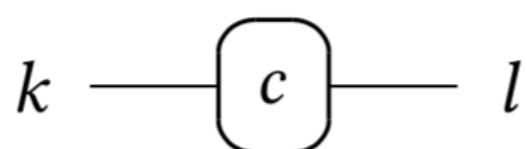
Stateless semantics can be inductively extended to a stateful one:

$$\llbracket c \rrbracket := \{(s_1, a, s_2, b) \mid (c, s_1) \xrightarrow[b]{a} (c, s_2)\}$$

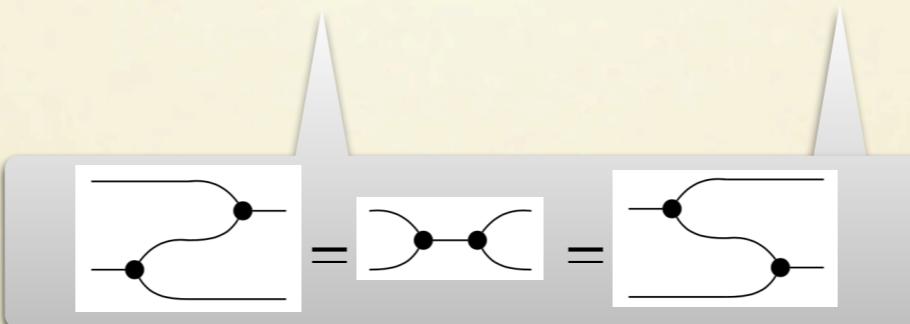
The register is canonical

Registers and their semantics are justified by the iso

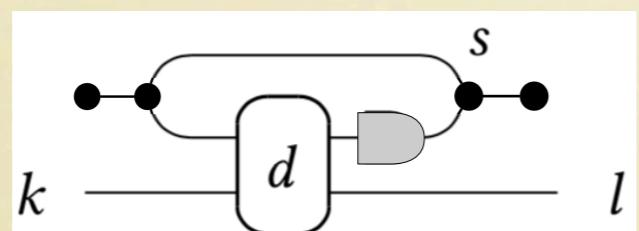
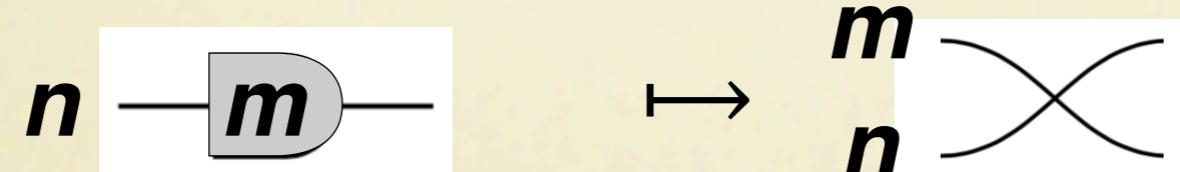
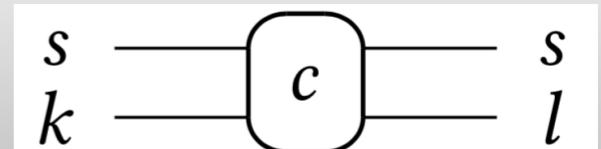
Arrows $k \rightarrow l$ are



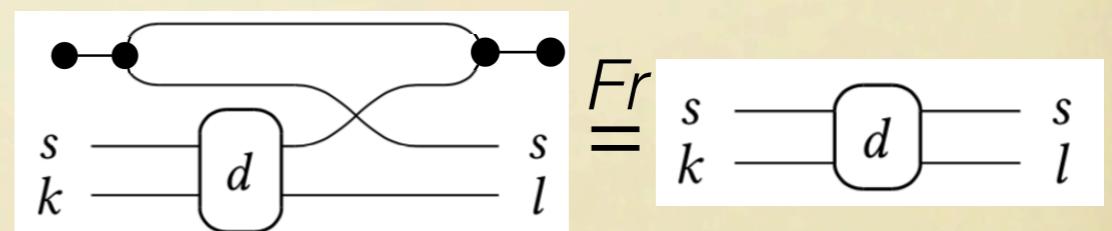
$$(\text{GLA}_S)_{/Fr} \cong \text{St}(\text{GLA}_{/Fr})$$



Arrows $k \rightarrow l$ are



→



What's coming up



The signal flow perspective

(CONCUR'14, POPL'15, *Inf. and Comp.* '17)

Linear interpretation

$$[[c]] := \{(a, b) \mid c \xrightarrow[b]{a} c\}$$

\mathbf{a} and \mathbf{b} are now vectors over \mathbb{R} .

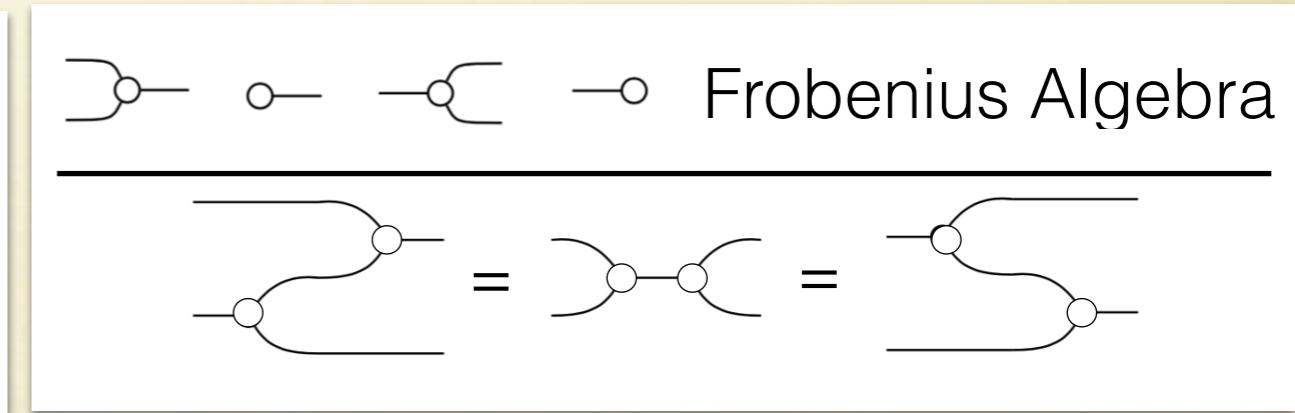
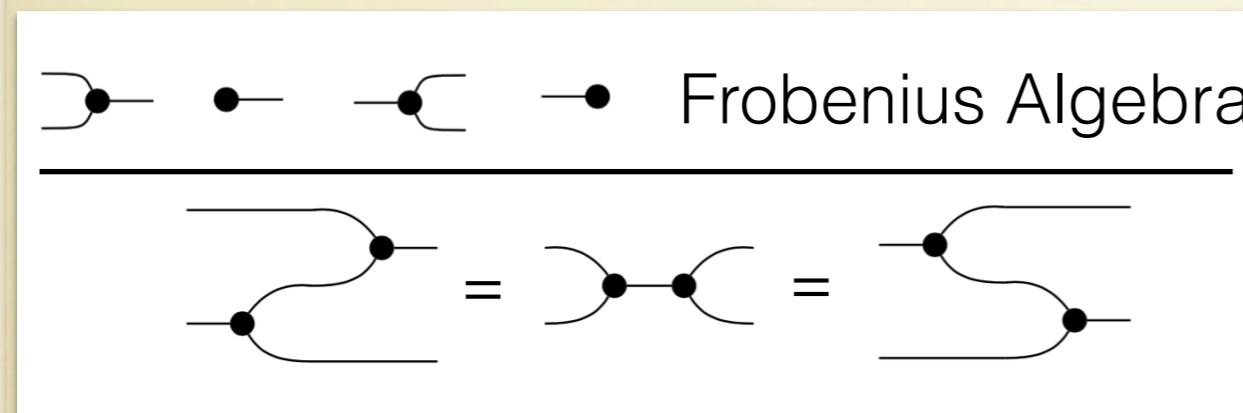
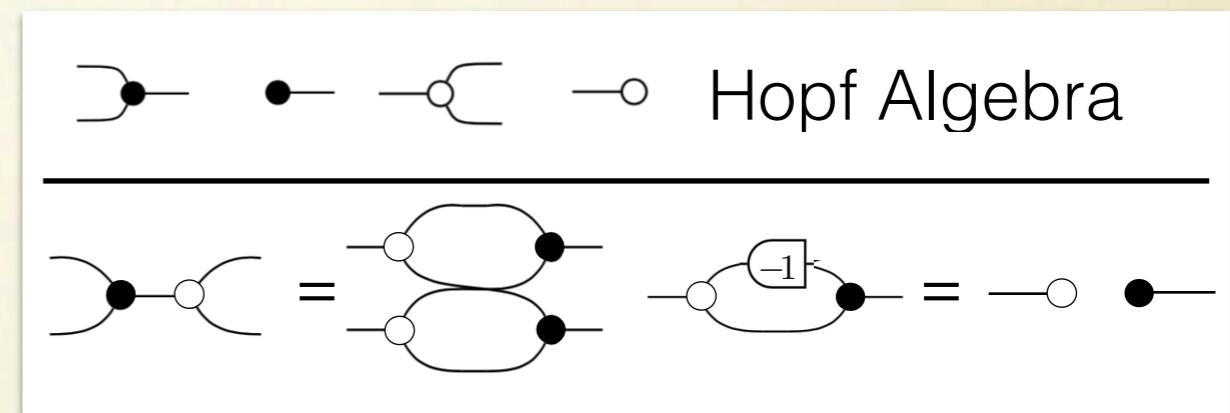
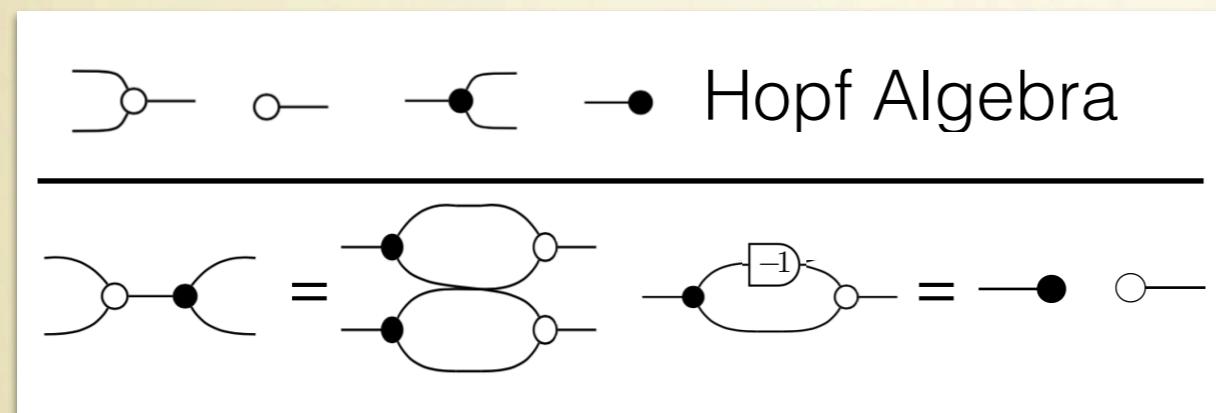
$[[c: k \rightarrow /]]$ is always a ***linear relation*** between \mathbb{R}^k and \mathbb{R}^l
(= a subspace of $\mathbb{R}^k \times \mathbb{R}^l$).

Proposition

Finite-dimensional linear relations form a category $\text{LinRel}_{\mathbb{R}}$

Equational Theory

III: Interacting Hopf Algebras



$$\square[n] \square[n] = \square[n] \square[n] \quad n \neq 0$$

Completeness

Theorem

$$\text{GLA}_{/\mathbb{H}} \cong \text{LinRel}_{\mathbb{R}}$$

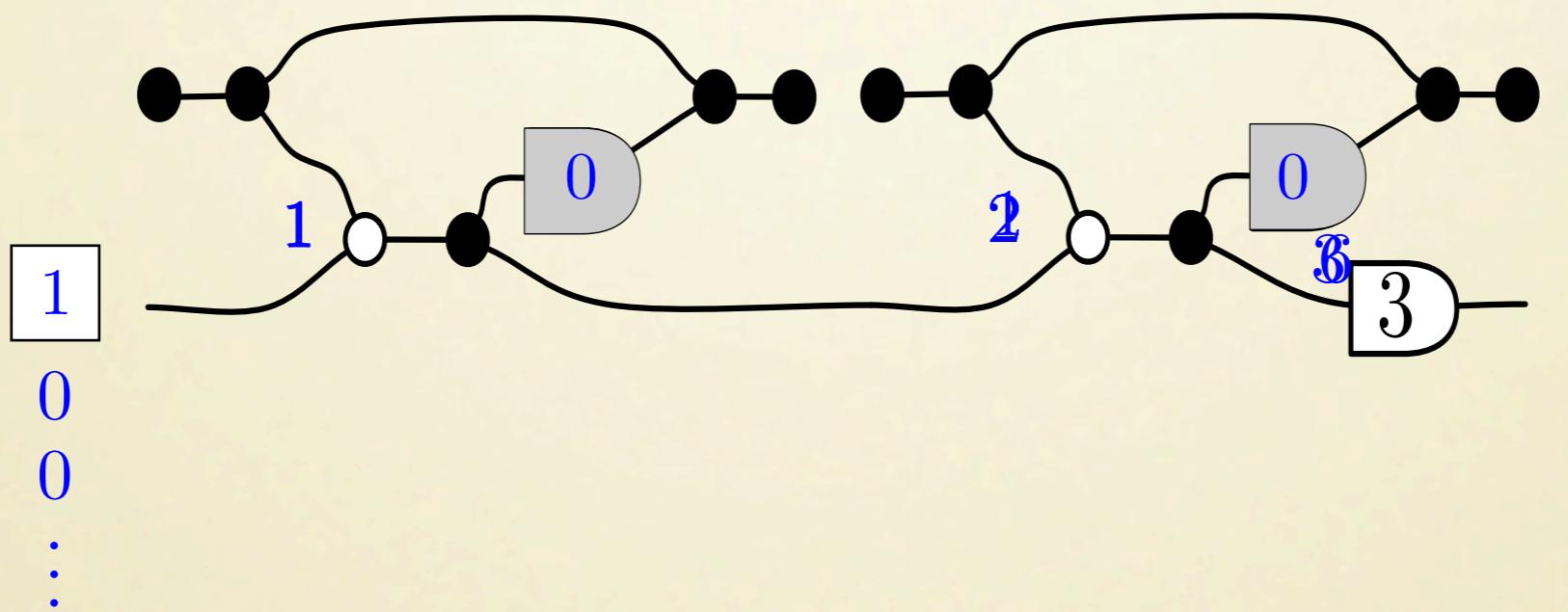
Corollary

$$[\![c]\!] = [\![d]\!] \iff c \stackrel{\mathbb{H}}{=} d$$

Corollary

$$(\text{GLA}_s)_{/\mathbb{H}} \cong \text{St}(\text{LinRel}_{\mathbb{R}})$$

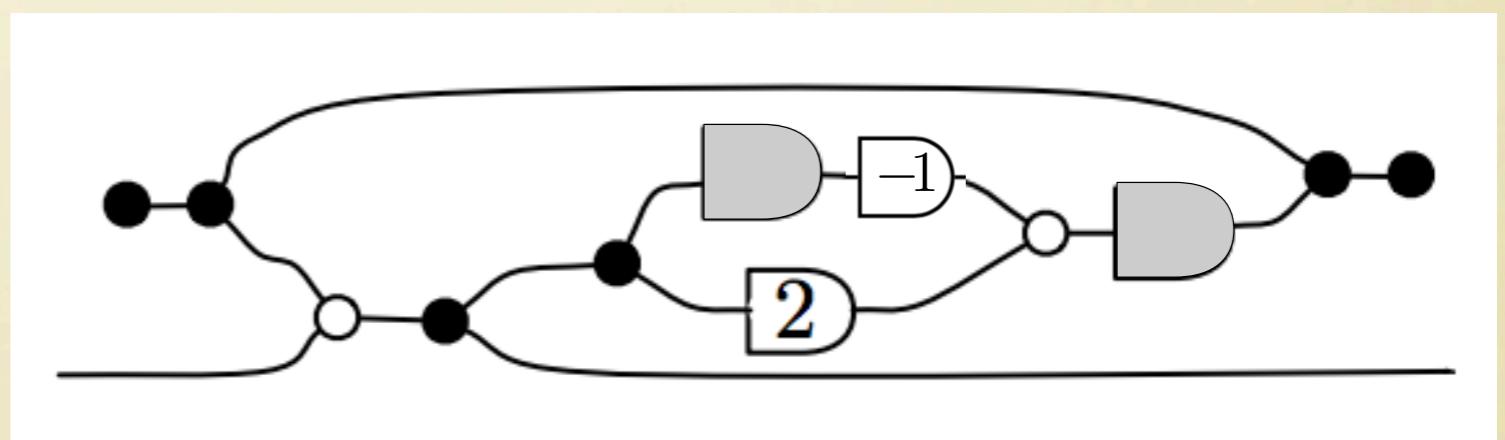
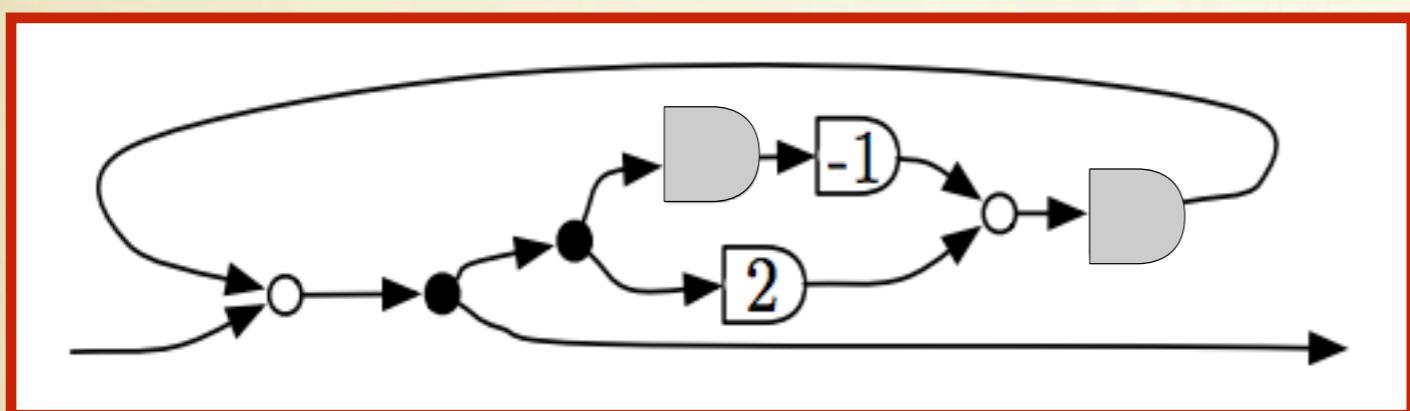
Stateful Example



$$(c,0,0) \xrightarrow[3]{1} (c,1,1) \xrightarrow[6]{0} (c,1,2) \xrightarrow[9]{0} (c,1,3) \xrightarrow[12]{0} \dots$$

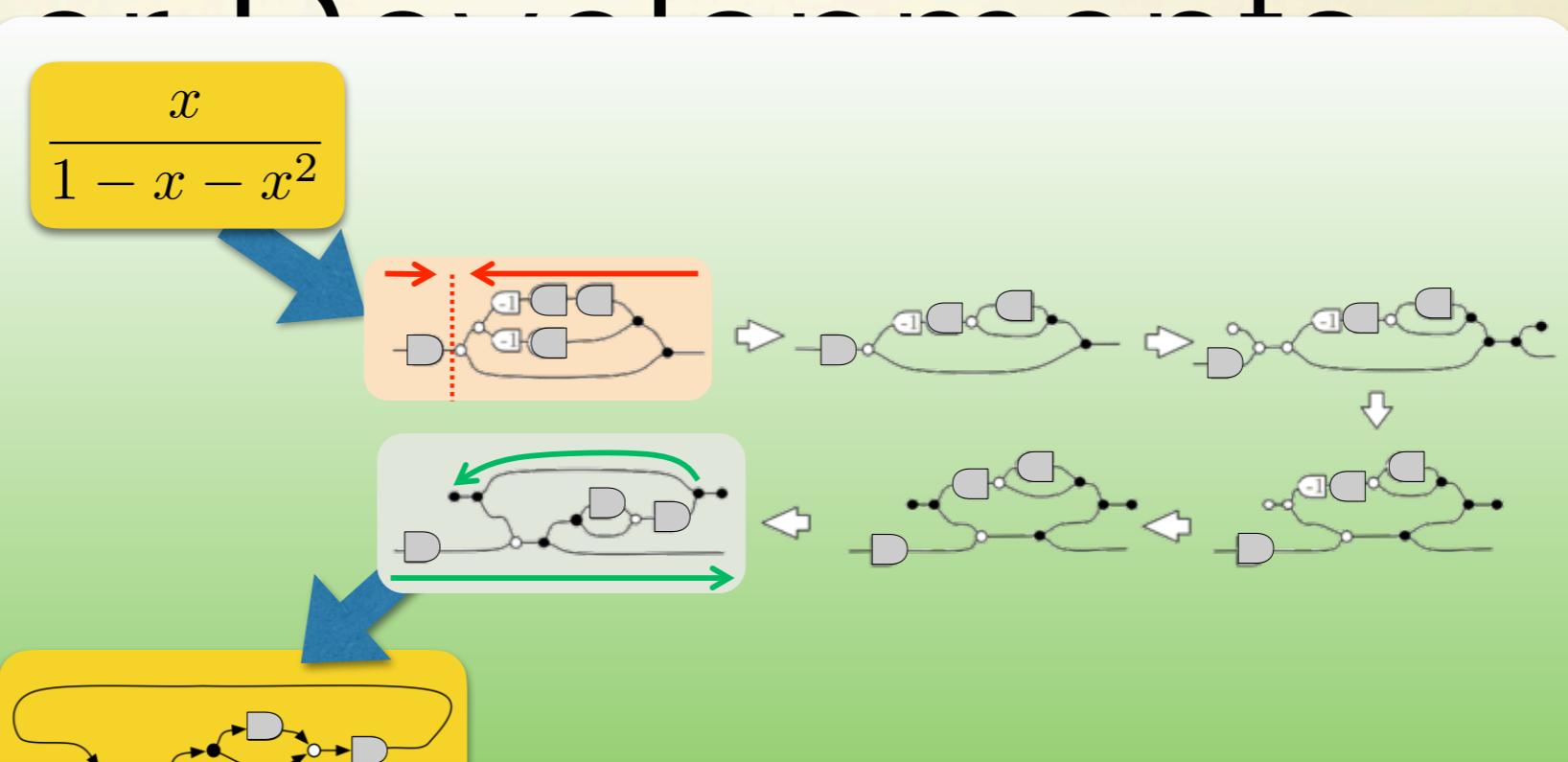
Signal Flow Graphs

In fact, the class of **signal flow graphs** embeds in GLA_s



Further Developments

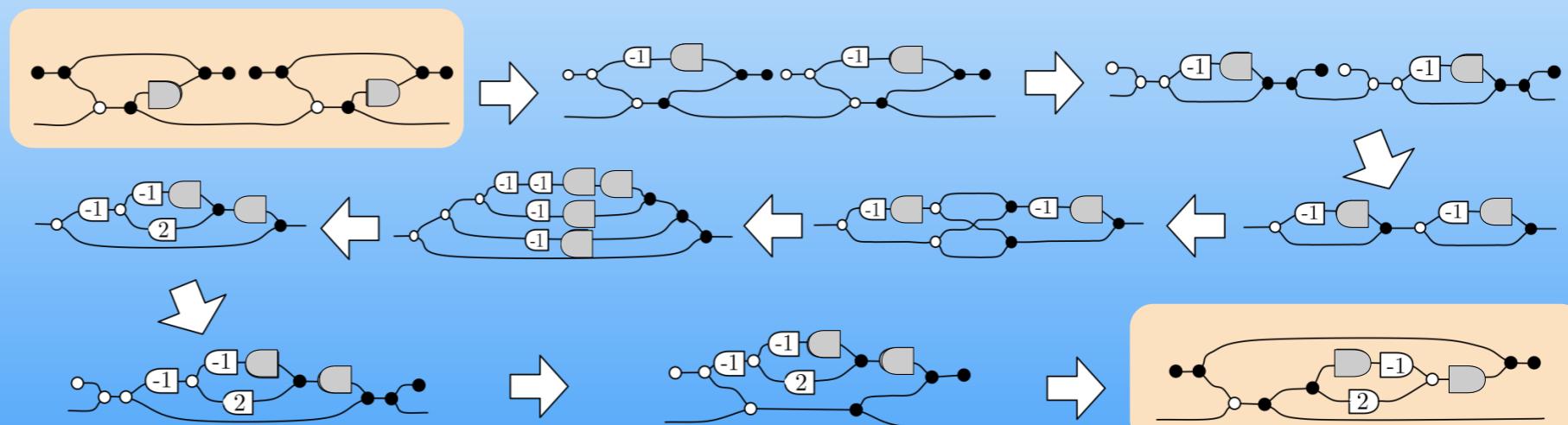
- ``Kleene's Theorem'' for Petri nets identify precisely the semantic domain.



- **Graphical Equations** for Petri Net Systems.
(Bonchi, Sobociński)



- **Realisability** problem
(Bonchi, Sobociński)

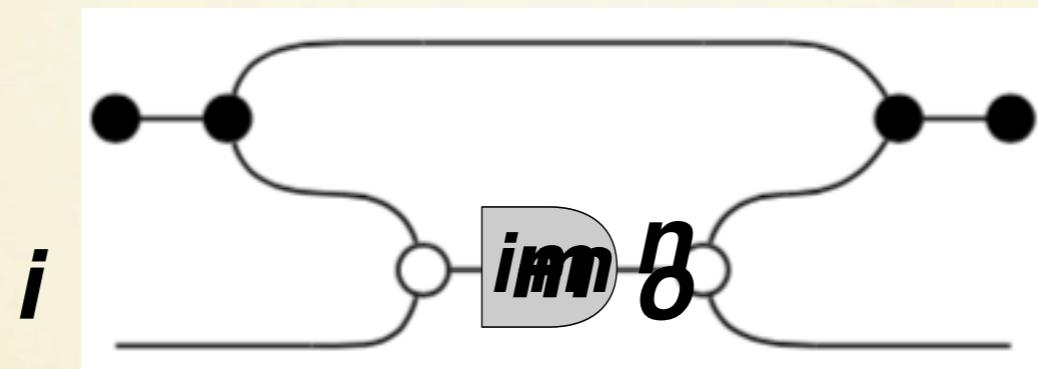


- Syntactic criteria
(Fong, Sobociński)

The resource perspective

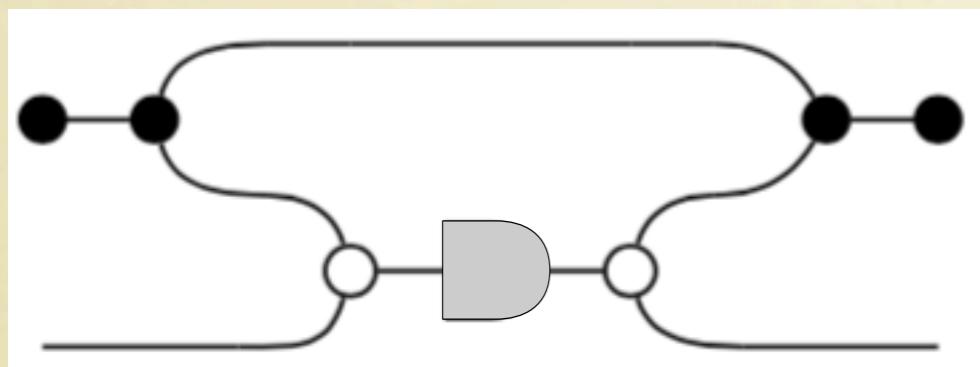
(POPL'19, LICS'19)

Motivating Example

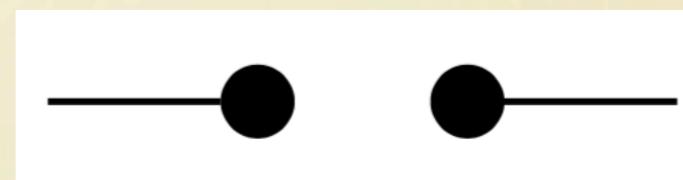


$$m = n + o$$

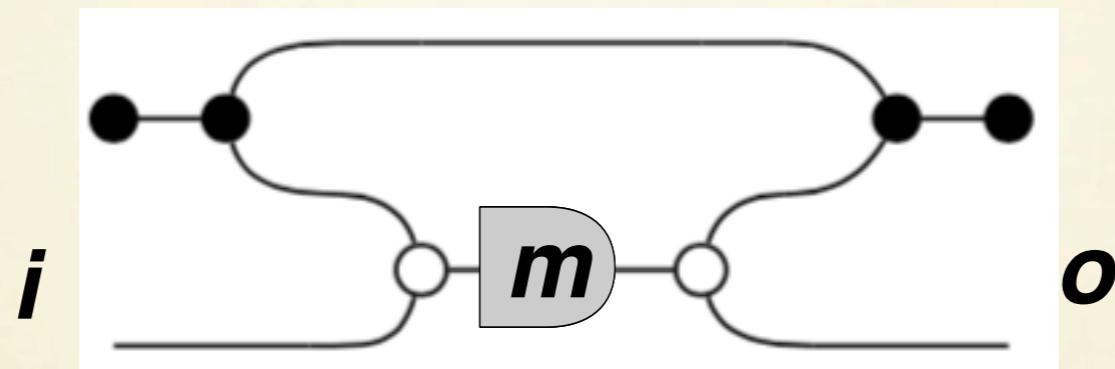
Over \mathbb{R} , any **i**, **o** and **m** will be in this relation.



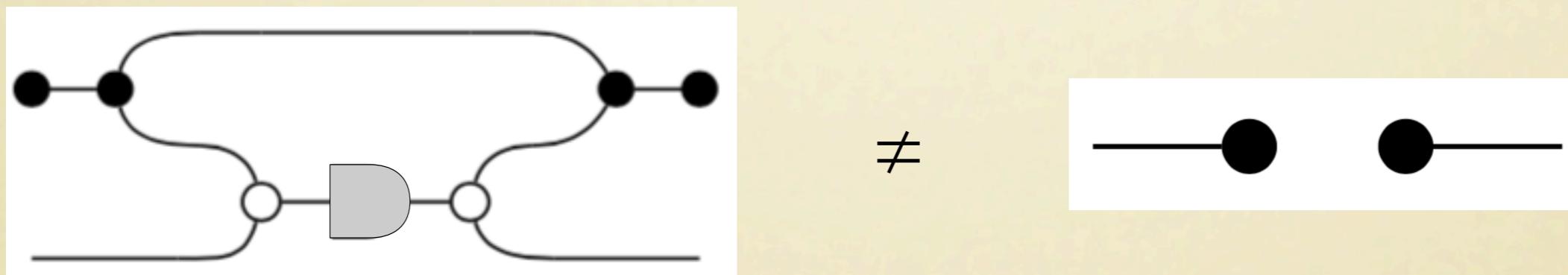
$$\mathbb{II} =$$



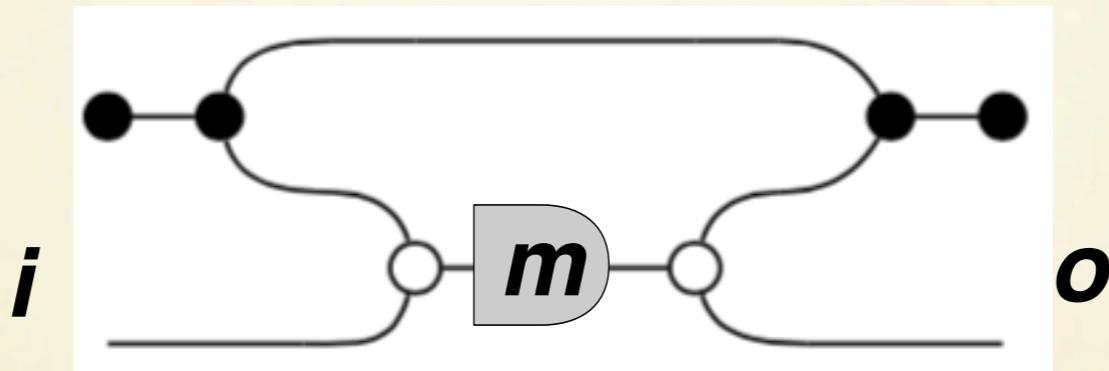
Motivating Example



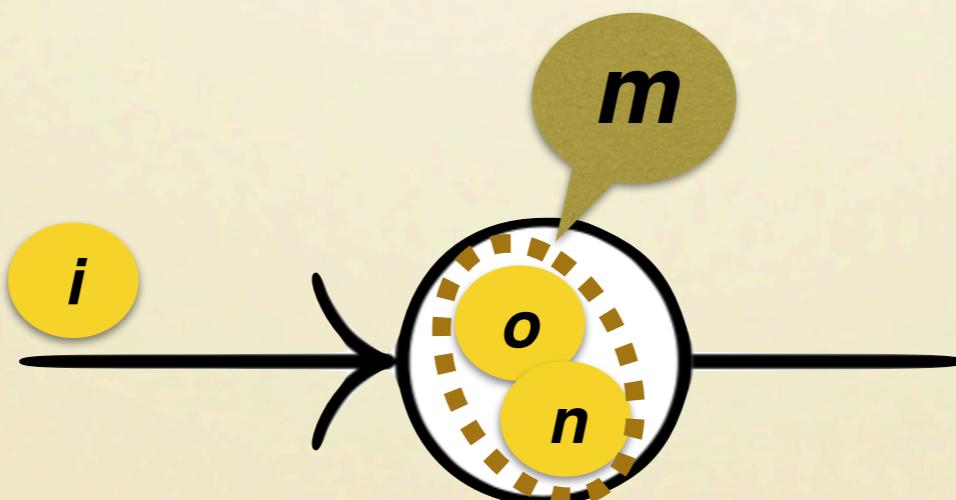
Over \mathbb{N} , ***i***, ***o*** and ***m*** are in the relation only if ***m* \geq *o***.



Motivating Example



In fact, over \mathbb{N} , this circuit behaves as the **place of a Petri net**.



Additive Relations

$$[[c]] := \{(a, b) \mid c \xrightarrow[b]{a} c\}$$

\mathbf{a} and \mathbf{b} are now vectors over \mathbb{N} .

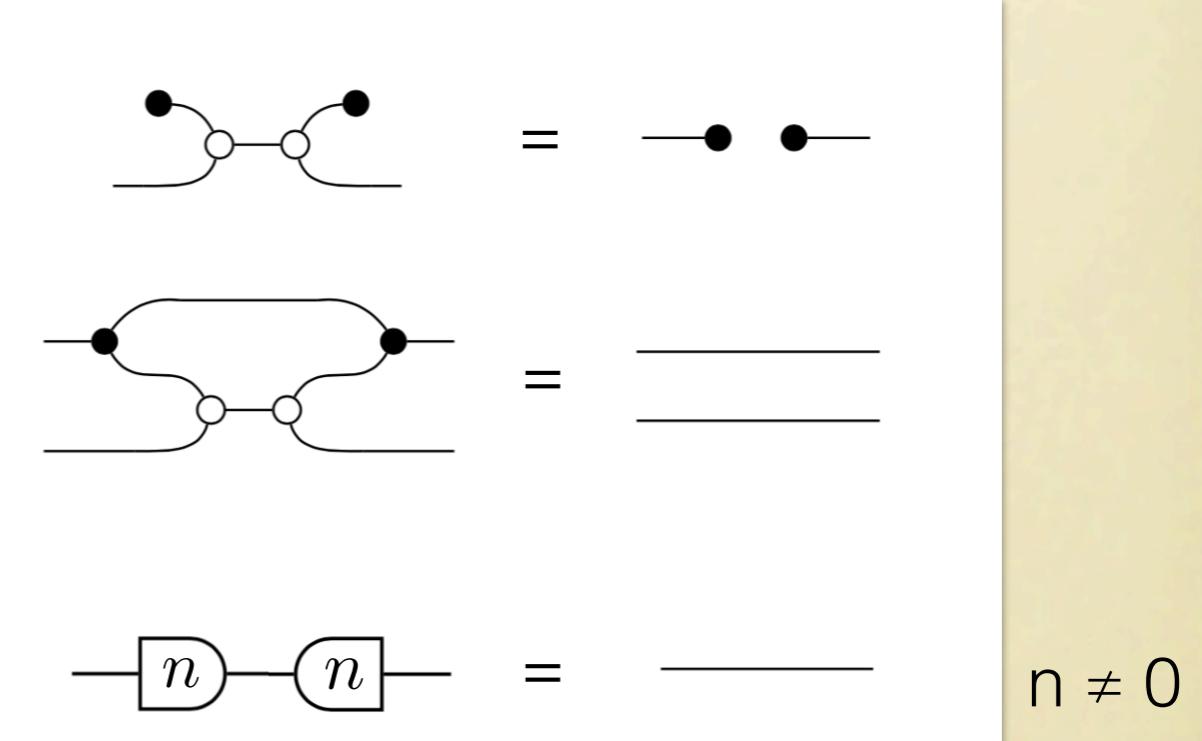
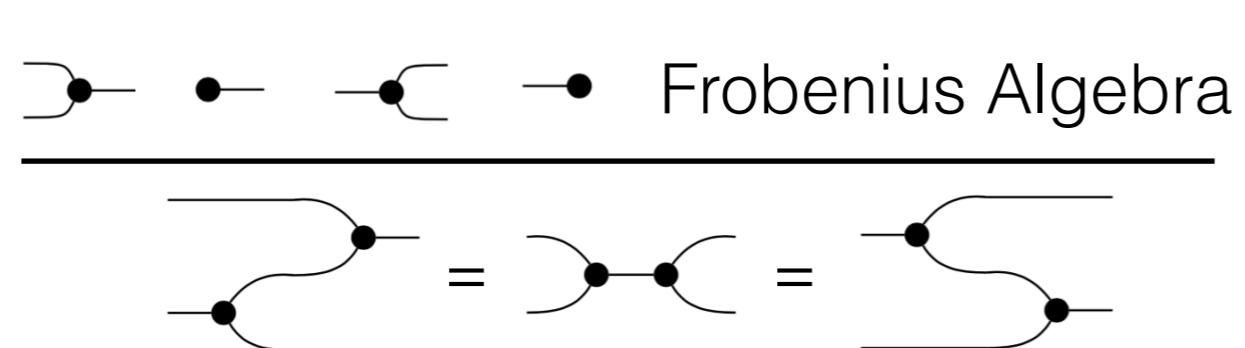
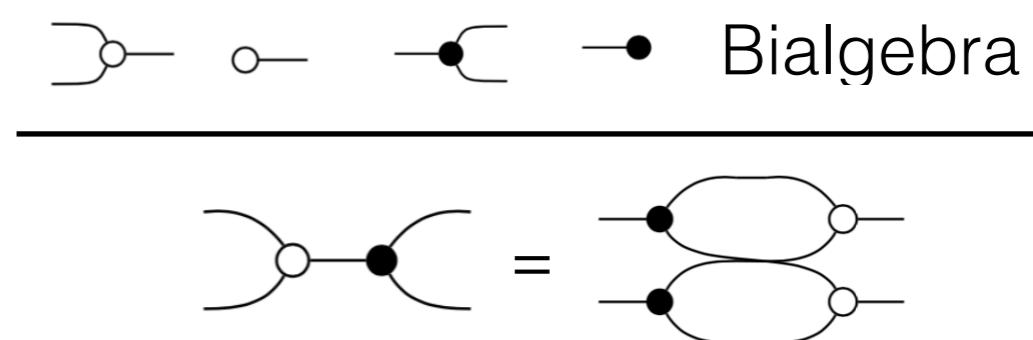
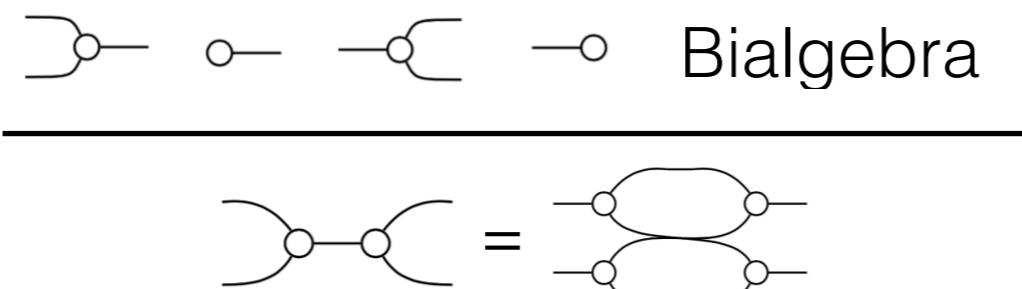
$[[c : k \rightarrow l]]$ is always an ***additive relation*** between \mathbb{N}^k and \mathbb{N}^l
(= a subset of $\mathbb{N}^k \times \mathbb{N}^l$ including $(\mathbf{0}, \mathbf{0})$ and closed under addition)

Proposition (non-trivial!)

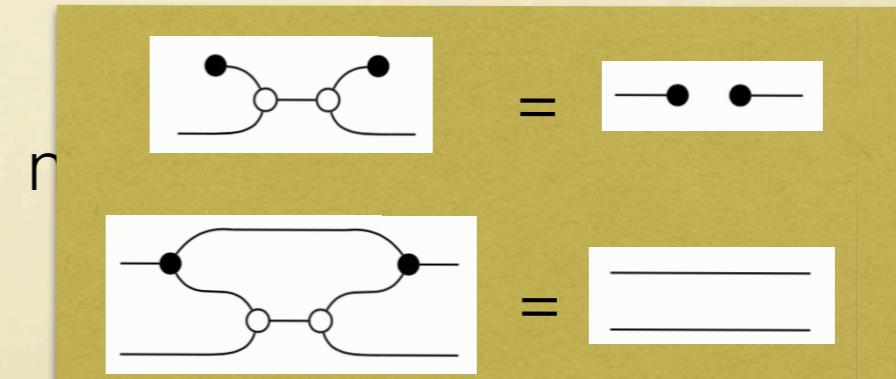
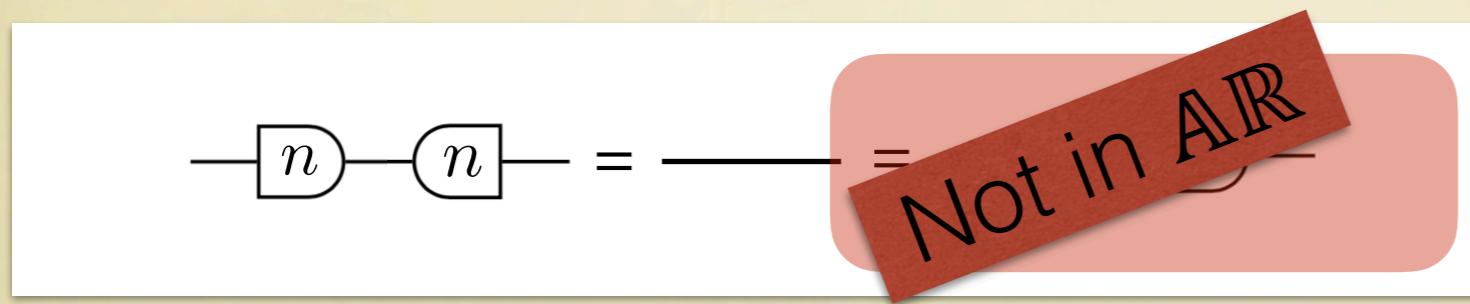
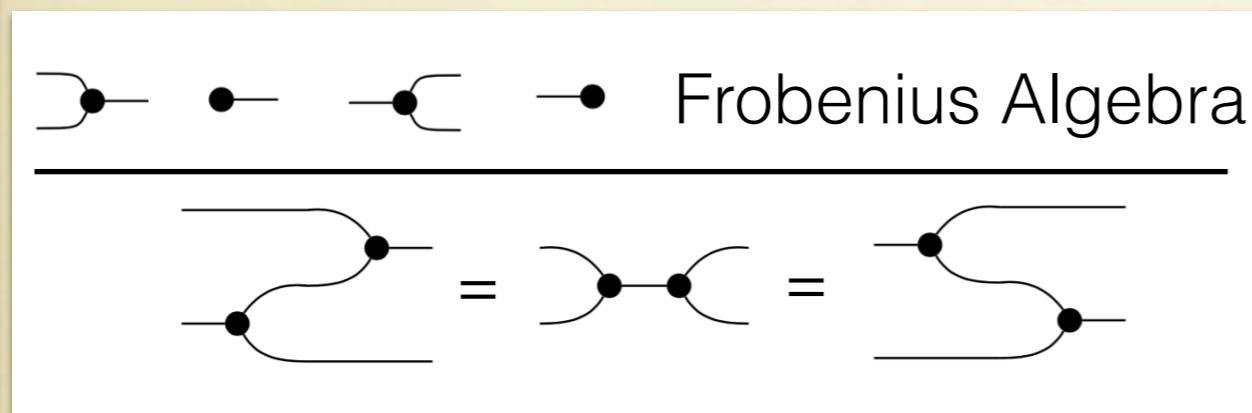
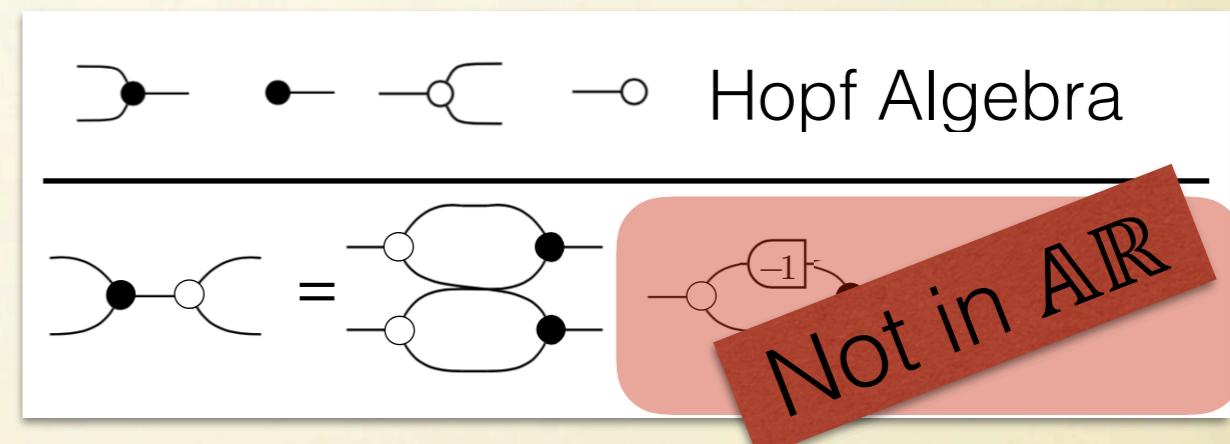
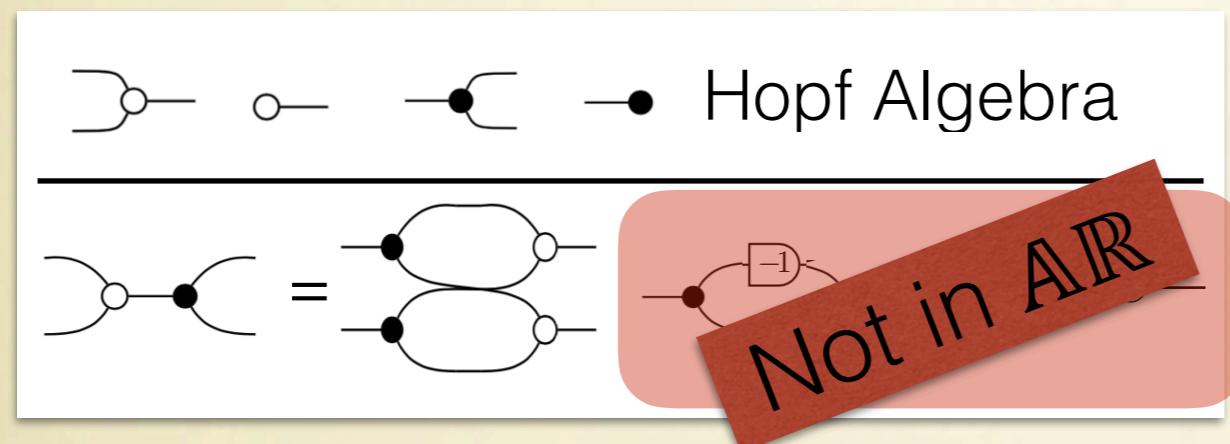
Finitely-generated additive relations form a category $\text{AddRel}_{\mathbb{N}}$

Equational Theory

AR: Algebra of Resources



AR vs IH



Completeness

Theorem

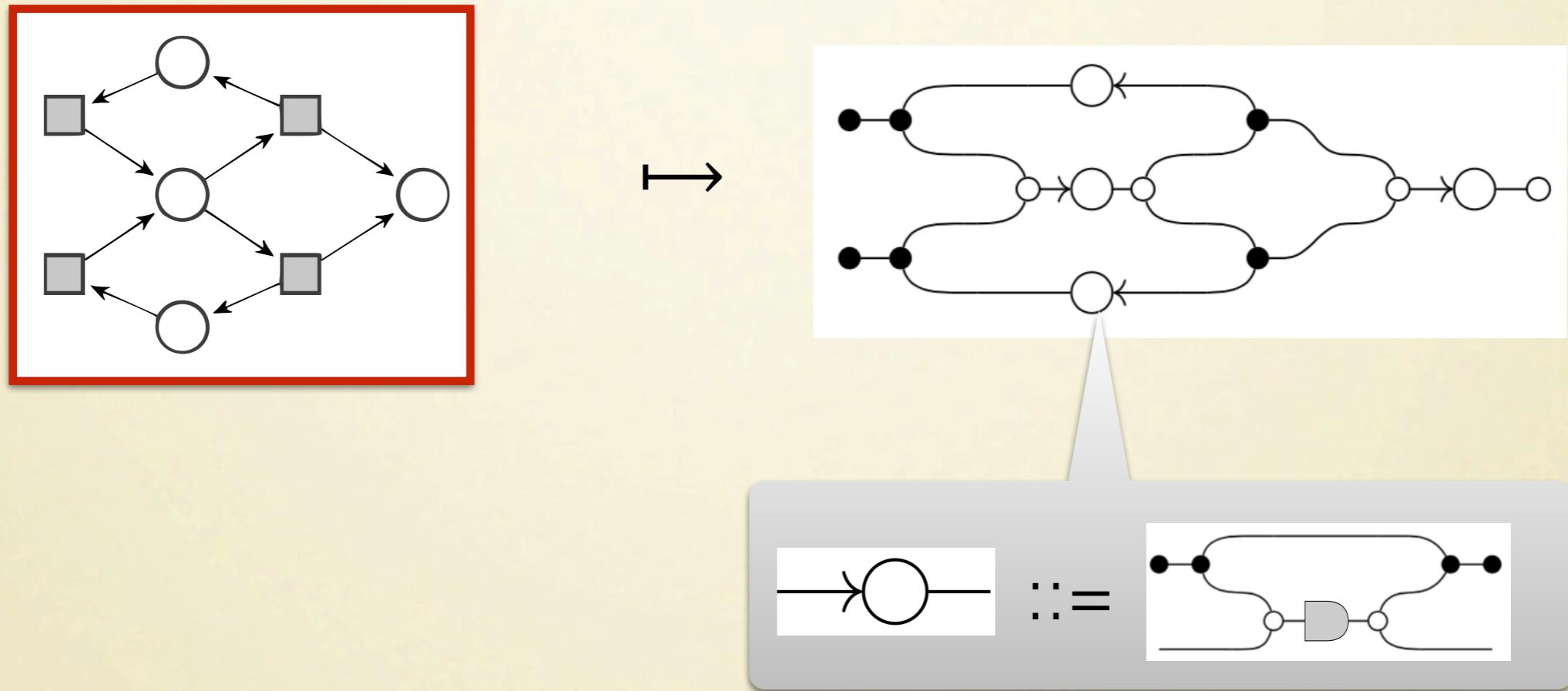
$$\text{GLA}_{/\mathbb{AR}} \cong \text{AddRel}_{\mathbb{N}}$$

$$[\![c]\!] = [\![d]\!] \iff c^{\mathbb{AR}} = d$$

Corollary

$$(\text{GLA}_s)_{/\mathbb{AR}} \cong \text{St}(\text{AddRel}_{\mathbb{N}})$$

Embedding Petri Nets



We can thus use \mathbb{AR} for equational reasoning about Petri Nets.

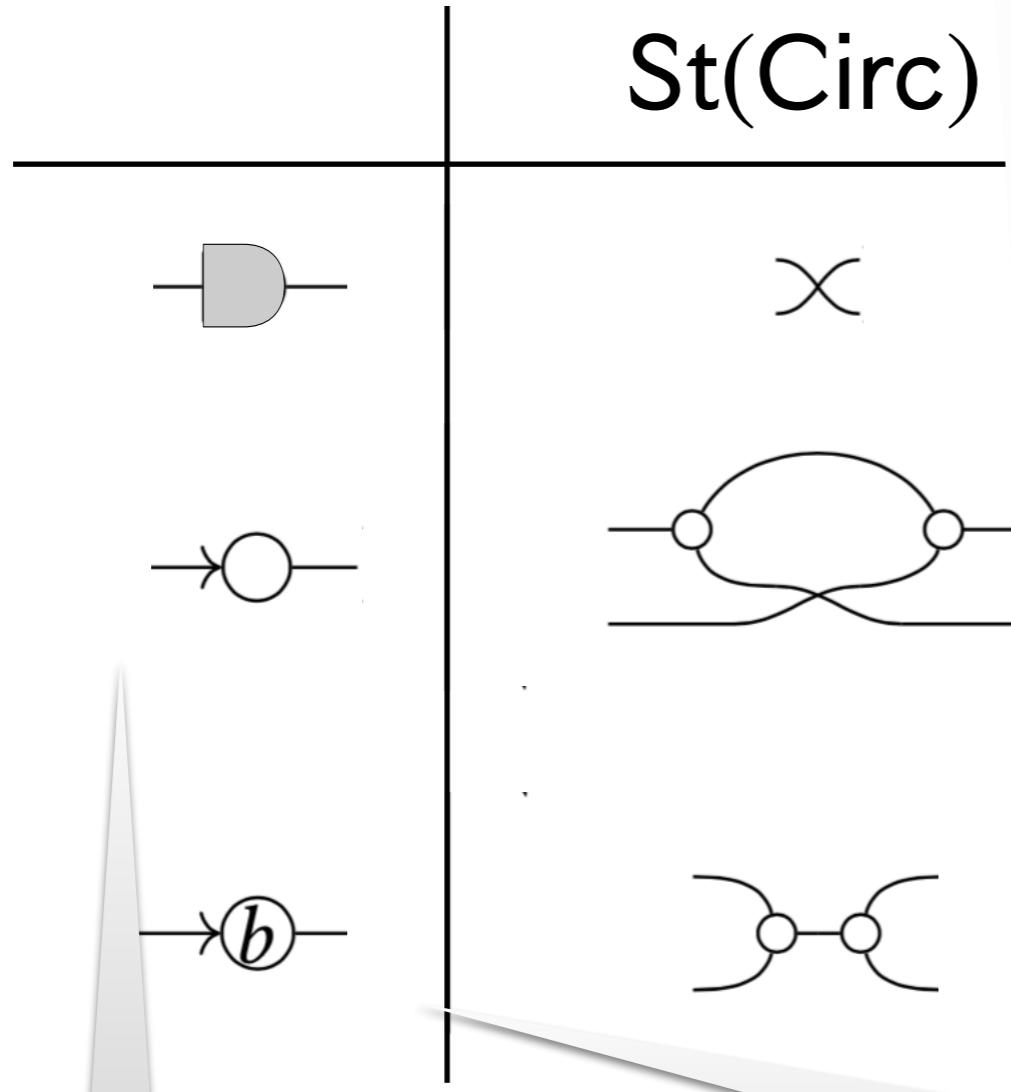
Classifying State Semantics

SEVIER

Electronic Notes in Theoretical Computer Science 162 (2006) 37–41

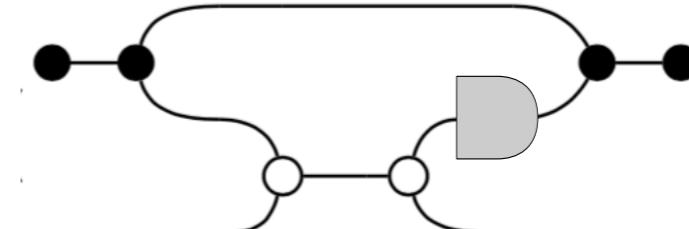
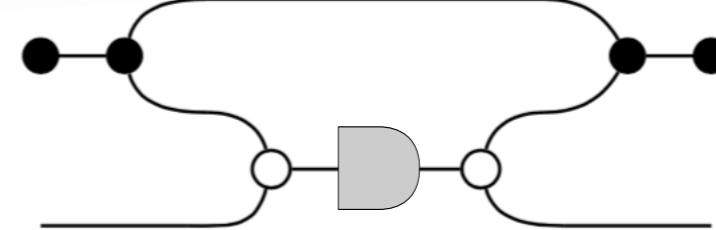
www.elsevier.com/locate/econline

St(Circ)



What are the Fundamental Structures of Concurrency?
We still don't know!

Samson Abramsky^{1,2}



Firing semantics

$$\frac{o \leq m}{(\rightarrow\circlearrowleft, m) \xrightarrow{o^i} (\rightarrow\circlearrowleft, m - o + i)}$$

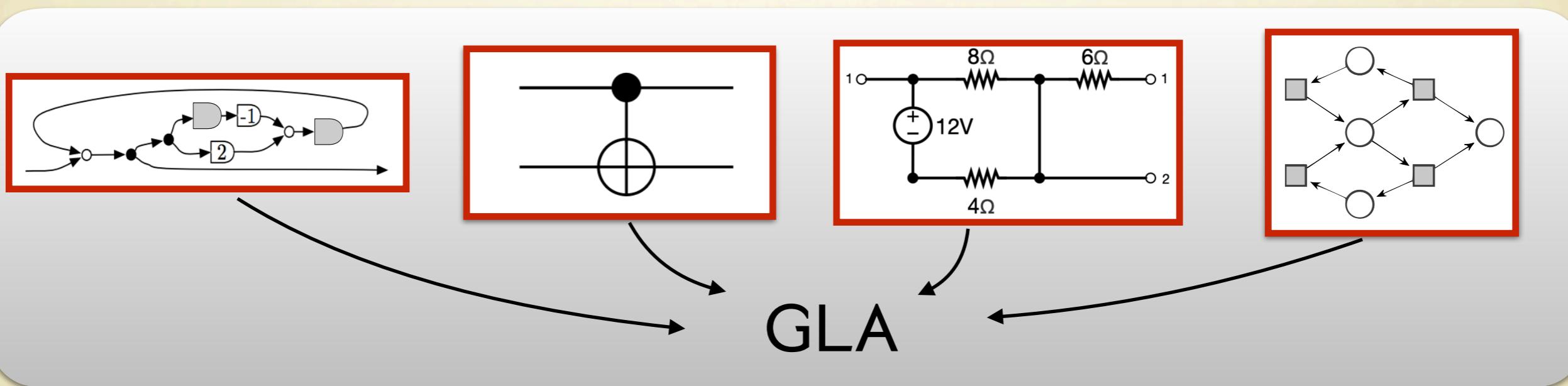
Banking semantics

$$\frac{m + i = m' + o}{(\rightarrow\circlearrowleft b, m) \xrightarrow{o^i} (\rightarrow\circlearrowleft b, m')}$$

Conclusions

Vision

We propose **GLA** as an **assembly language** for diverse families of component-based systems.



Major research thread: cover **probabilistic** models
(Bayesian reasoning, machine learning, etc.).

Bibliography

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